SNU 4541.574
Programming Language Theory

Ack: BCP's slides

More on Types

Review: Typing Rules

$$\texttt{true} : \texttt{Bool} \qquad \text{(T-True)}$$

$$\texttt{false} : \texttt{Bool} \qquad \text{(T-False)}$$

$$\frac{\texttt{t}_1 : \texttt{Bool} \quad \texttt{t}_2 : \texttt{T} \quad \texttt{t}_3 : \texttt{T}}{\texttt{if } \texttt{t}_1 \texttt{ then } \texttt{t}_2 \texttt{ else } \texttt{t}_3 : \texttt{T}} \qquad \text{(T-If)}$$

$$\texttt{0} : \texttt{Nat} \qquad \text{(T-Zero)}$$

$$\frac{\texttt{t}_1 : \texttt{Nat}}{\texttt{succ } \texttt{t}_1 : \texttt{Nat}} \qquad \text{(T-Succ)}$$

$$\frac{\texttt{t}_1 : \texttt{Nat}}{\texttt{pred } \texttt{t}_1 : \texttt{Nat}} \qquad \text{(T-Pred)}$$

$$\frac{\texttt{t}_1 : \texttt{Nat}}{\texttt{iszero } \texttt{t}_1 : \texttt{Bool}} \qquad \text{(T-IsZero)}$$

Review: Inversion

Lemma:

1. If $\texttt{true} : \texttt{R}$, then $\texttt{R} = \texttt{Bool}$.
2. If $\texttt{false} : \texttt{R}$, then $\texttt{R} = \texttt{Bool}$.
3. If $\texttt{if } \texttt{t}_1 \texttt{ then } \texttt{t}_2 \texttt{ else } \texttt{t}_3 : \texttt{R}$, then $\texttt{t}_1 : \texttt{Bool}$, $\texttt{t}_2 : \texttt{R}$, and $\texttt{t}_3 : \texttt{R}$.
4. If $\texttt{0} : \texttt{R}$, then $\texttt{R} = \texttt{Nat}$.
5. If $\texttt{succ } \texttt{t}_1 : \texttt{R}$, then $\texttt{R} = \texttt{Nat}$ and $\texttt{t}_1 : \texttt{Nat}$.
6. If $\texttt{pred } \texttt{t}_1 : \texttt{R}$, then $\texttt{R} = \texttt{Nat}$ and $\texttt{t}_1 : \texttt{Nat}$.
7. If $\texttt{iszero } \texttt{t}_1 : \texttt{R}$, then $\texttt{R} = \texttt{Bool}$ and $\texttt{t}_1 : \texttt{Nat}$.

Canonical Forms

Lemma:

1. If `v` is a value of type `Bool`, then `v` is either `true` or `false`.
2. If `v` is a value of type `Nat`, then `v` is a numeric value.

Proof:

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Proof: Recall the syntax of values:

`v` ::=		*values*
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`nv` ::=		*numeric values*
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For part 1, if `v` is `true` or `false`, the result is immediate.

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For part 1, if `v` is `true` or `false`, the result is immediate. But `v` cannot be `0` or `succ nv`, since the inversion lemma tells us that `v` would then have type `Nat`, not `Bool`. Part 2 is similar.

Progress

Theorem: Suppose $\mathtt{t}$ is a well-typed term (that is, $\mathtt{t} : \mathtt{T}$ for some type $\mathtt{T}$). Then either $\mathtt{t}$ is a value or else there is some $\mathtt{t}'$ with $\mathtt{t} \longrightarrow \mathtt{t}'$.

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Theorem: Suppose $\texttt{t}$ is a well-typed term (that is, $\texttt{t : T}$ for some type $\texttt{T}$). Then either $\texttt{t}$ is a value or else there is some $\texttt{t}'$ with $\texttt{t} \longrightarrow \texttt{t}'$.

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The T-True, T-False, and T-Zero cases are immediate, since $\mathtt{t}$ in these cases is a value.

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Case T-If: $\mathtt{t} = \mathtt{if}\ \mathtt{t}_1\ \mathtt{then}\ \mathtt{t}_2\ \mathtt{else}\ \mathtt{t}_3$

$\mathtt{t}_1 : \mathtt{Bool} \qquad \mathtt{t}_2 : \mathtt{T} \qquad \mathtt{t}_3 : \mathtt{T}$

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The T-True, T-False, and T-Zero cases are immediate, since $\mathtt{t}$ in these cases is a value.

Case T-If: $\mathtt{t} = \mathtt{if\ t_1\ then\ t_2\ else\ t_3}$

$\mathtt{t_1 : Bool} \qquad \mathtt{t_2 : T} \qquad \mathtt{t_3 : T}$

By the induction hypothesis, either $\mathtt{t_1}$ is a value or else there is some $\mathtt{t_1'}$ such that $\mathtt{t_1} \longrightarrow \mathtt{t_1'}$. If $\mathtt{t_1}$ is a value, then the canonical forms lemma tells us that it must be either $\mathtt{true}$ or $\mathtt{false}$, in which case either E-IfTrue or E-IfFalse applies to $\mathtt{t}$. On the other hand, if $\mathtt{t_1} \longrightarrow \mathtt{t_1'}$, then, by E-If, $\mathtt{t} \longrightarrow \mathtt{if\ t_1'\ then\ t_2\ else\ t_3}$.

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Proof: By induction on a derivation of $\mathtt{t : T}$.

The cases for rules T-Zero, T-Succ, T-Pred, and T-IsZero are similar.

(Recommended: Try to reconstruct them.)

Preservation

Theorem: If $\texttt{t} : \texttt{T}$ and $\texttt{t} \longrightarrow \texttt{t}'$, then $\texttt{t}' : \texttt{T}$.

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Case T-True: $\quad \mathtt{t} = \mathtt{true} \qquad \mathtt{T} = \mathtt{Bool}$

Then $\mathtt{t}$ is a value, so it cannot be that $\mathtt{t} \longrightarrow \mathtt{t}'$ for any $\mathtt{t}'$, and the theorem is vacuously true.

Preservation

Theorem: If $\mathtt{t} : \mathtt{T}$ and $\mathtt{t} \longrightarrow \mathtt{t}'$, then $\mathtt{t}' : \mathtt{T}$.

Proof: By induction on the given typing derivation.

Case T-If:

$$\mathtt{t} = \mathtt{if}\ \mathtt{t}_1\ \mathtt{then}\ \mathtt{t}_2\ \mathtt{else}\ \mathtt{t}_3 \quad \mathtt{t}_1 : \mathtt{Bool} \quad \mathtt{t}_2 : \mathtt{T} \quad \mathtt{t}_3 : \mathtt{T}$$

There are three evaluation rules by which $\mathtt{t} \longrightarrow \mathtt{t}'$ can be derived: E-IfTrue, E-IfFalse, and E-If. Consider each case separately.

Preservation

Theorem: If $\mathtt{t} : \mathtt{T}$ and $\mathtt{t} \longrightarrow \mathtt{t}'$, then $\mathtt{t}' : \mathtt{T}$.

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There are three evaluation rules by which $\mathtt{t} \longrightarrow \mathtt{t}'$ can be derived: E-IfTrue, E-IfFalse, and E-If. Consider each case separately.

Subcase E-IfTrue: $\quad \mathtt{t}_1 = \mathtt{true} \qquad \mathtt{t}' = \mathtt{t}_2$

Immediate, by the assumption $\mathtt{t}_2 : \mathtt{T}$.

(E-IfFalse subcase: Similar.)

Preservation

Theorem: If $\mathtt{t} : \mathtt{T}$ and $\mathtt{t} \longrightarrow \mathtt{t}'$, then $\mathtt{t}' : \mathtt{T}$.

Proof: By induction on the given typing derivation.

Case T-IF:

$$\mathtt{t} = \mathtt{if}\ \mathtt{t}_1\ \mathtt{then}\ \mathtt{t}_2\ \mathtt{else}\ \mathtt{t}_3 \quad \mathtt{t}_1 : \mathtt{Bool} \quad \mathtt{t}_2 : \mathtt{T} \quad \mathtt{t}_3 : \mathtt{T}$$

There are three evaluation rules by which $\mathtt{t} \longrightarrow \mathtt{t}'$ can be derived: E-IFTRUE, E-IFFALSE, and E-IF. Consider each case separately.

Subcase E-IF: $\mathtt{t}_1 \longrightarrow \mathtt{t}_1'$ $\quad \mathtt{t}' = \mathtt{if}\ \mathtt{t}_1'\ \mathtt{then}\ \mathtt{t}_2\ \mathtt{else}\ \mathtt{t}_3$

Applying the IH to the subderivation of $\mathtt{t}_1 : \mathtt{Bool}$ yields $\mathtt{t}_1' : \mathtt{Bool}$. Combining this with the assumptions that $\mathtt{t}_2 : \mathtt{T}$ and $\mathtt{t}_3 : \mathtt{T}$, we can apply rule T-IF to conclude that $\mathtt{if}\ \mathtt{t}_1'\ \mathtt{then}\ \mathtt{t}_2\ \mathtt{else}\ \mathtt{t}_3 : \mathtt{T}$, that is, $\mathtt{t}' : \mathtt{T}$.

The Simply Typed Lambda-Calculus

The simply typed lambda-calculus

The system we are about to define is commonly called the *simply typed lambda-calculus*, or $\lambda_{\rightarrow}$ for short.

Unlike the untyped lambda-calculus, the "pure" form of $\lambda_{\rightarrow}$ (with no primitive values or operations) is not very interesting; to talk about $\lambda_{\rightarrow}$, we always begin with some set of "base types."

- So, strictly speaking, there are *many* variants of $\lambda_{\rightarrow}$, depending on the choice of base types.
- For now, we'll work with a variant constructed over the booleans.

Untyped lambda-calculus with booleans

$\texttt{t}$	::=		*terms*
		$\texttt{x}$	*variable*
		$\lambda\texttt{x.t}$	*abstraction*
		$\texttt{t t}$	*application*
		$\texttt{true}$	*constant true*
		$\texttt{false}$	*constant false*
		$\texttt{if t then t else t}$	*conditional*
$\texttt{v}$	::=		*values*
		$\lambda\texttt{x.t}$	*abstraction value*
		$\texttt{true}$	*true value*
		$\texttt{false}$	*false value*

"Simple Types"

`T` ::=		*types*
	`Bool`	*type of booleans*
	`T→T`	*types of functions*

Type Annotations

We now have a choice to make. Do we...

- annotate lambda-abstractions with the expected type of the argument

  $$\lambda \texttt{x:T}_1\texttt{. t}_2$$

  (as in most mainstream programming languages), or
- continue to write lambda-abstractions as before

  $$\lambda \texttt{x. t}_2$$

  and ask the typing rules to “guess” an appropriate annotation (as in OCaml)?

Both are reasonable choices, but the first makes the job of defining the typin rules simpler. Let’s take this choice for now.

Typing rules

$$\texttt{true} : \texttt{Bool} \qquad \text{(T-True)}$$

$$\texttt{false} : \texttt{Bool} \qquad \text{(T-False)}$$

$$\frac{\texttt{t}_1 : \texttt{Bool} \qquad \texttt{t}_2 : \texttt{T} \qquad \texttt{t}_3 : \texttt{T}}{\texttt{if } \texttt{t}_1 \texttt{ then } \texttt{t}_2 \texttt{ else } \texttt{t}_3 : \texttt{T}} \qquad \text{(T-If)}$$

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$$\frac{???}{\lambda \texttt{x}:\texttt{T}_1.\texttt{t}_2 : \texttt{T}_1 \to \texttt{T}_2} \qquad \text{(T-Abs)}$$

Typing rules

$$\texttt{true} : \texttt{Bool} \qquad \text{(T-TRUE)}$$

$$\texttt{false} : \texttt{Bool} \qquad \text{(T-FALSE)}$$

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$$\frac{\Gamma, \texttt{x}:\texttt{T}_1 \vdash \texttt{t}_2 : \texttt{T}_2}{\Gamma \vdash \lambda \texttt{x}:\texttt{T}_1 . \texttt{t}_2 : \texttt{T}_1 \to \texttt{T}_2} \qquad \text{(T-ABS)}$$

$$\frac{\texttt{x}:\texttt{T} \in \Gamma}{\Gamma \vdash \texttt{x} : \texttt{T}} \qquad \text{(T-VAR)}$$

Typing rules

$$\Gamma \vdash \texttt{true} : \texttt{Bool} \qquad \text{(T-True)}$$

$$\Gamma \vdash \texttt{false} : \texttt{Bool} \qquad \text{(T-False)}$$

$$\frac{\Gamma \vdash \texttt{t}_1 : \texttt{Bool} \quad \Gamma \vdash \texttt{t}_2 : \texttt{T} \quad \Gamma \vdash \texttt{t}_3 : \texttt{T}}{\Gamma \vdash \texttt{if } \texttt{t}_1 \texttt{ then } \texttt{t}_2 \texttt{ else } \texttt{t}_3 : \texttt{T}} \qquad \text{(T-If)}$$

$$\frac{\Gamma, \texttt{x}{:}\texttt{T}_1 \vdash \texttt{t}_2 : \texttt{T}_2}{\Gamma \vdash \lambda \texttt{x}{:}\texttt{T}_1 . \texttt{t}_2 : \texttt{T}_1 \to \texttt{T}_2} \qquad \text{(T-Abs)}$$

$$\frac{\texttt{x}{:}\texttt{T} \in \Gamma}{\Gamma \vdash \texttt{x} : \texttt{T}} \qquad \text{(T-Var)}$$

$$\frac{\Gamma \vdash \texttt{t}_1 : \texttt{T}_{11} \to \texttt{T}_{12} \quad \Gamma \vdash \texttt{t}_2 : \texttt{T}_{11}}{\Gamma \vdash \texttt{t}_1 \; \texttt{t}_2 : \texttt{T}_{12}} \qquad \text{(T-App)}$$

Typing Derivations

What derivations justify the following typing statements?

- $\vdash$ `(λx:Bool.x) true` : `Bool`
- `f:Bool→Bool` $\vdash$ `f (if false then true else false)` : `Bool`
- `f:Bool→Bool` $\vdash$ `λx:Bool. f (if x then false else x)` : `Bool→Bool`

Properties of $\lambda_{\rightarrow}$

The fundamental property of the type system we have just defined is *soundness* with respect to the operational semantics.

1. *Progress:* A closed, well-typed term is not stuck

   If $\vdash t : T$*, then either* t *is a value or else* $t \longrightarrow t'$ *for some* t'*.*

2. *Preservation:* Types are preserved by one-step evaluation

   If $\Gamma \vdash t : T$ *and* $t \longrightarrow t'$*, then* $\Gamma \vdash t' : T$*.*

Proving progress

Same steps as before...

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- inversion lemma for typing relation
- canonical forms lemma
- progress theorem

Inversion

Lemma:

1. If $\Gamma \vdash \texttt{true} : \texttt{R}$, then $\texttt{R} = \texttt{Bool}$.
2. If $\Gamma \vdash \texttt{false} : \texttt{R}$, then $\texttt{R} = \texttt{Bool}$.
3. If $\Gamma \vdash \texttt{if } \texttt{t}_1 \texttt{ then } \texttt{t}_2 \texttt{ else } \texttt{t}_3 : \texttt{R}$, then $\Gamma \vdash \texttt{t}_1 : \texttt{Bool}$ and $\Gamma \vdash \texttt{t}_2, \texttt{t}_3 : \texttt{R}$.

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4. If $\Gamma \vdash \texttt{x} : \texttt{R}$, then $\texttt{x:R} \in \Gamma$.

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5. If $\Gamma \vdash \lambda \mathtt{x}{:}\mathtt{T}_1.\mathtt{t}_2 : \mathtt{R}$, then $\mathtt{R} = \mathtt{T}_1 {\to} \mathtt{R}_2$ for some $\mathtt{R}_2$ with $\Gamma, \mathtt{x}{:}\mathtt{T}_1 \vdash \mathtt{t}_2 : \mathtt{R}_2$.

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5. If $\Gamma \vdash \lambda\texttt{x}{:}\texttt{T}_1\texttt{.t}_2 : \texttt{R}$, then $\texttt{R} = \texttt{T}_1 {\to} \texttt{R}_2$ for some $\texttt{R}_2$ with $\Gamma, \texttt{x}{:}\texttt{T}_1 \vdash \texttt{t}_2 : \texttt{R}_2$.
6. If $\Gamma \vdash \texttt{t}_1\ \texttt{t}_2 : \texttt{R}$, then there is some type $\texttt{T}_{11}$ such that $\Gamma \vdash \texttt{t}_1 : \texttt{T}_{11} {\to} \texttt{R}$ and $\Gamma \vdash \texttt{t}_2 : \texttt{T}_{11}$.

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2. If v is a value of type $T_1 \to T_2$, then v has the form $\lambda x{:}T_1.t_2$.

Progress

Theorem: Suppose $\mathtt{t}$ is a closed, well-typed term (that is, $\vdash \mathtt{t} : \mathtt{T}$ for some $\mathtt{T}$). Then either $\mathtt{t}$ is a value or else there is some $\mathtt{t}'$ with $\mathtt{t} \longrightarrow \mathtt{t}'$.

Proof: By induction

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Proof: By induction on typing derivations. The cases for boolean constants and conditions are the same as before. The variable case is trivial (because $\mathtt{t}$ is closed). The abstraction case is immediate, since abstractions are values.

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Proof: By induction on typing derivations. The cases for boolean constants and conditions are the same as before. The variable case is trivial (because $\mathtt{t}$ is closed). The abstraction case is immediate, since abstractions are values.
Consider the case for application, where $\mathtt{t} = \mathtt{t}_1\ \mathtt{t}_2$ with $\vdash \mathtt{t}_1 : \mathtt{T}_{11} \rightarrow \mathtt{T}_{12}$ and $\vdash \mathtt{t}_2 : \mathtt{T}_{11}$.

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Theorem: Suppose $\mathtt{t}$ is a closed, well-typed term (that is, $\vdash \mathtt{t} : \mathtt{T}$ for some $\mathtt{T}$). Then either $\mathtt{t}$ is a value or else there is some $\mathtt{t}'$ with $\mathtt{t} \longrightarrow \mathtt{t}'$.

Proof: By induction on typing derivations. The cases for boolean constants and conditions are the same as before. The variable case is trivial (because $\mathtt{t}$ is closed). The abstraction case is immediate, since abstractions are values.
Consider the case for application, where $\mathtt{t} = \mathtt{t}_1\ \mathtt{t}_2$ with $\vdash \mathtt{t}_1 : \mathtt{T}_{11} \rightarrow \mathtt{T}_{12}$ and $\vdash \mathtt{t}_2 : \mathtt{T}_{11}$. By the induction hypothesis, either $\mathtt{t}_1$ is a value or else it can make a step of evaluation, and likewise $\mathtt{t}_2$.

Progress

Theorem: Suppose $\texttt{t}$ is a closed, well-typed term (that is, $\vdash \texttt{t} : \texttt{T}$ for some $\texttt{T}$). Then either $\texttt{t}$ is a value or else there is some $\texttt{t}'$ with $\texttt{t} \longrightarrow \texttt{t}'$.

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Consider the case for application, where $\texttt{t} = \texttt{t}_1\ \texttt{t}_2$ with $\vdash \texttt{t}_1 : \texttt{T}_{11}\to\texttt{T}_{12}$ and $\vdash \texttt{t}_2 : \texttt{T}_{11}$. By the induction hypothesis, either $\texttt{t}_1$ is a value or else it can make a step of evaluation, and likewise $\texttt{t}_2$. If $\texttt{t}_1$ can take a step, then rule E-App1 applies to $\texttt{t}$. If $\texttt{t}_1$ is a value and $\texttt{t}_2$ can take a step, then rule E-App2 applies. Finally, if both $\texttt{t}_1$ and $\texttt{t}_2$ are values, then the canonical forms lemma tells us that $\texttt{t}_1$ has the form $\lambda\texttt{x}:\texttt{T}_{11}.\texttt{t}_{12}$, and so rule E-AppAbs applies to $\texttt{t}$.

Preservation

Theorem: If $\Gamma \vdash \mathtt{t} : \mathtt{T}$ and $\mathtt{t} \longrightarrow \mathtt{t}'$, then $\Gamma \vdash \mathtt{t}' : \mathtt{T}$.

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Which case is the hard one??

Preservation

Theorem: If $\Gamma \vdash \mathtt{t} : \mathtt{T}$ and $\mathtt{t} \longrightarrow \mathtt{t}'$, then $\Gamma \vdash \mathtt{t}' : \mathtt{T}$.

Proof: By induction on typing derivations.

Case T-APP: Given $\mathtt{t} = \mathtt{t}_1\ \mathtt{t}_2$

$\Gamma \vdash \mathtt{t}_1 : \mathtt{T}_{11} \rightarrow \mathtt{T}_{12}$

$\Gamma \vdash \mathtt{t}_2 : \mathtt{T}_{11}$

$\mathtt{T} = \mathtt{T}_{12}$

Show $\Gamma \vdash \mathtt{t}' : \mathtt{T}_{12}$

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$\mathtt{T} = \mathtt{T}_{12}$

Show $\Gamma \vdash \mathtt{t}' : \mathtt{T}_{12}$

By the inversion lemma for evaluation, there are three subcases...

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Subcase: $\mathtt{t}_1 = \lambda \mathtt{x}{:}\mathtt{T}_{11}.\ \mathtt{t}_{12}$

$\mathtt{t}_2$ a value $\mathtt{v}_2$

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Uh oh.

The “Substitution Lemma”

Lemma: Types are preserved under substititition.

That is, if $\Gamma, \mathtt{x}{:}\mathtt{S} \vdash \mathtt{t} : \mathtt{T}$ and $\Gamma \vdash \mathtt{s} : \mathtt{S}$, then $\Gamma \vdash [\mathtt{x} \mapsto \mathtt{s}]\mathtt{t} : \mathtt{T}$.

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Proof: ...

Preservation

Recommended: Complete the proof of preservation