

# Corpus based approach in Constraint solving

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# Reference

- S. Gulwani, S. Srivastava, and R. Venkatesan,  
Program analysis as constraint solving, PLDI'08

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# Motivation

- How is the corpus-based (or enumeration based) approach used well in constraint solving?

# Background

- Discovering program invariant
  - fixed-point computation based approach (e.g. abstract interpretation)
  - constraint-based invariant generation approach
    - program  $\longrightarrow$  satisfiability constraints

# Background

- Two advantages of Constraint-based approach
  - more efficient
  - more precise

# Constraint solving

- Applications
  - program verification
  - strongest postcondition generation
  - weakest precondition generation

# Program verification

- Goal
  - verifying assertions in program are valid



# Program verification

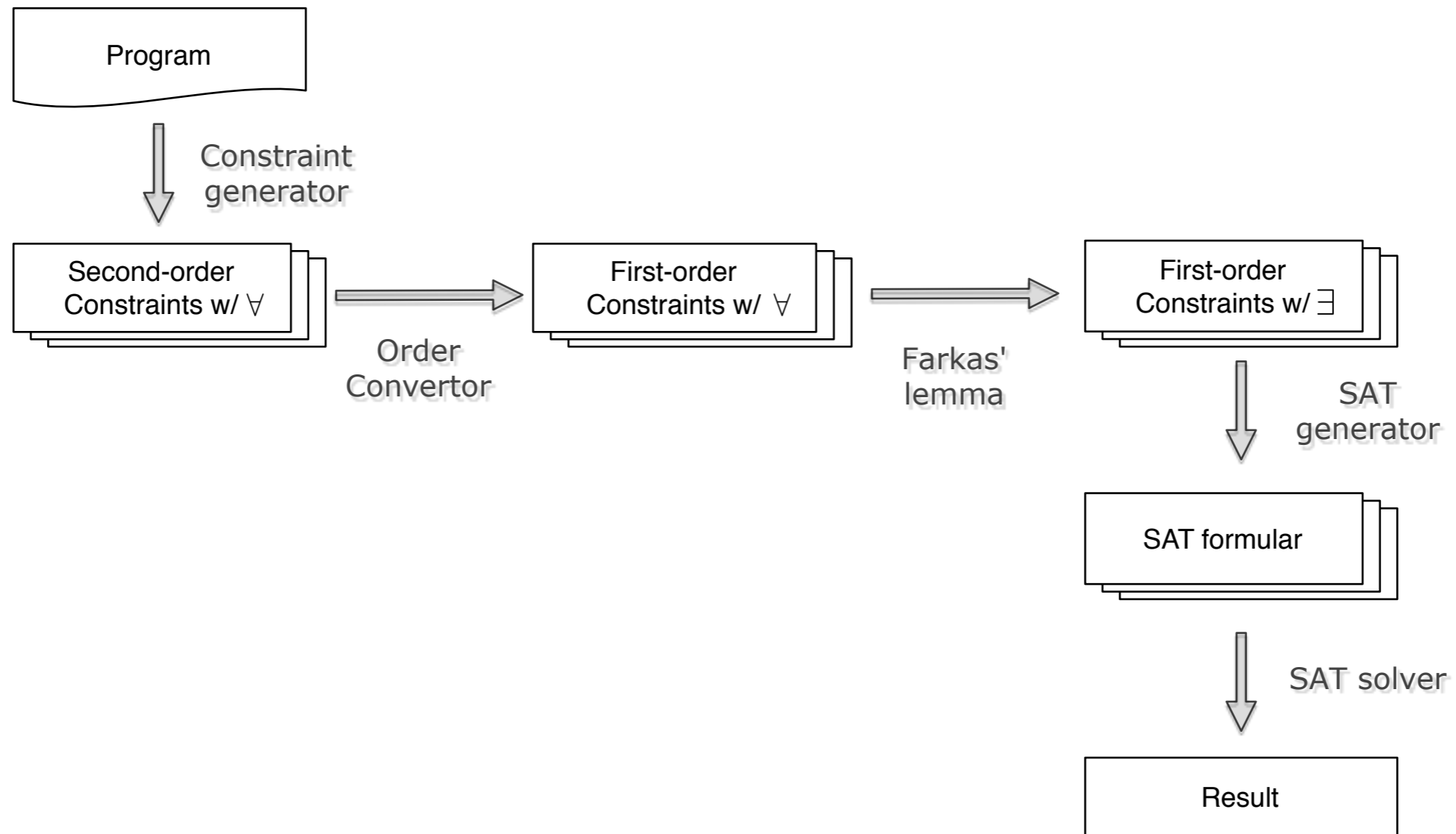
```
PV1 (int y) {  
    x := -50;  
    while (x < 0) {  
        x := x + y;  
        y++;  
    }  
    assert(y > 0)  
}
```

# Program verification

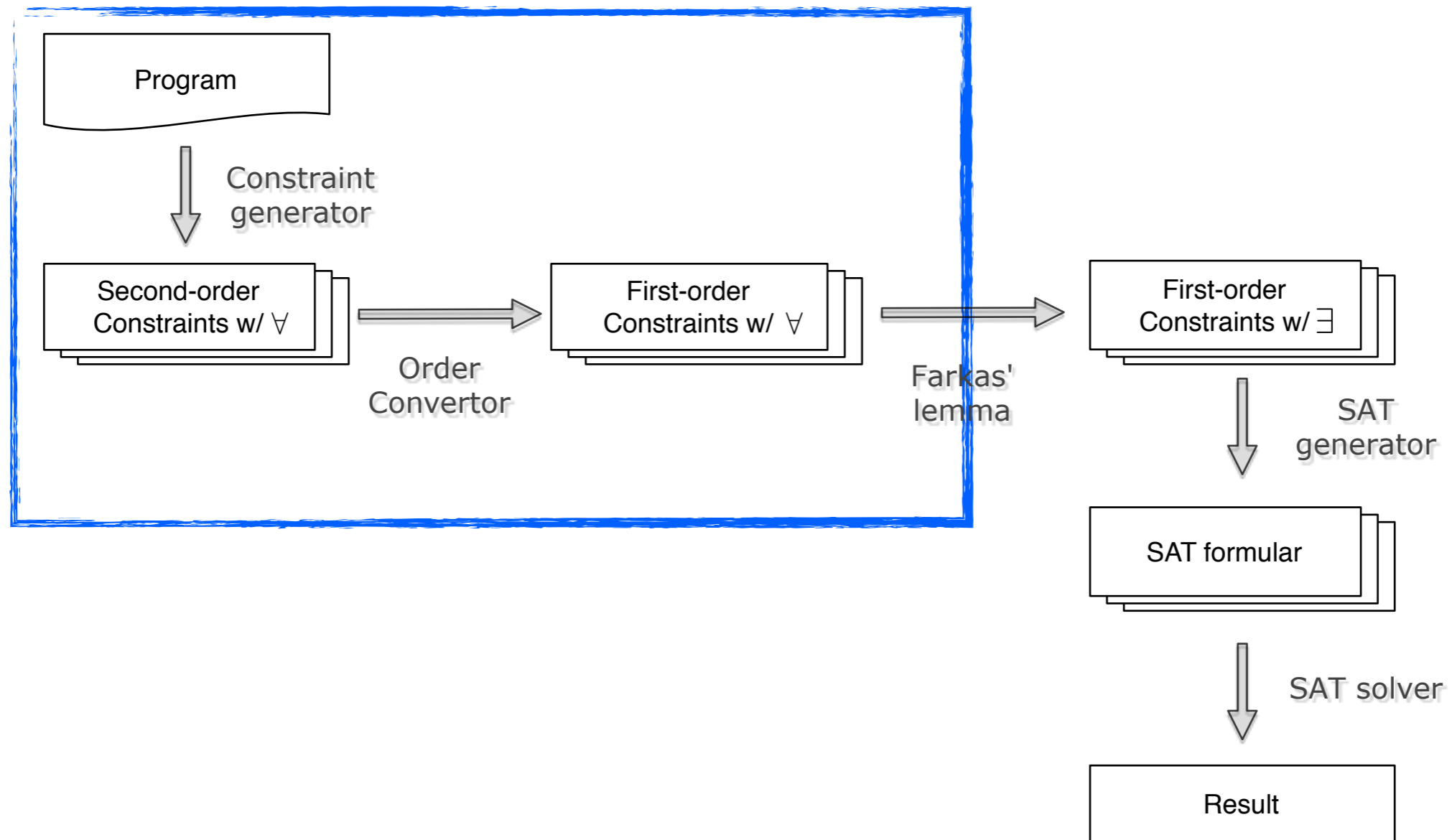
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Valid or Not?

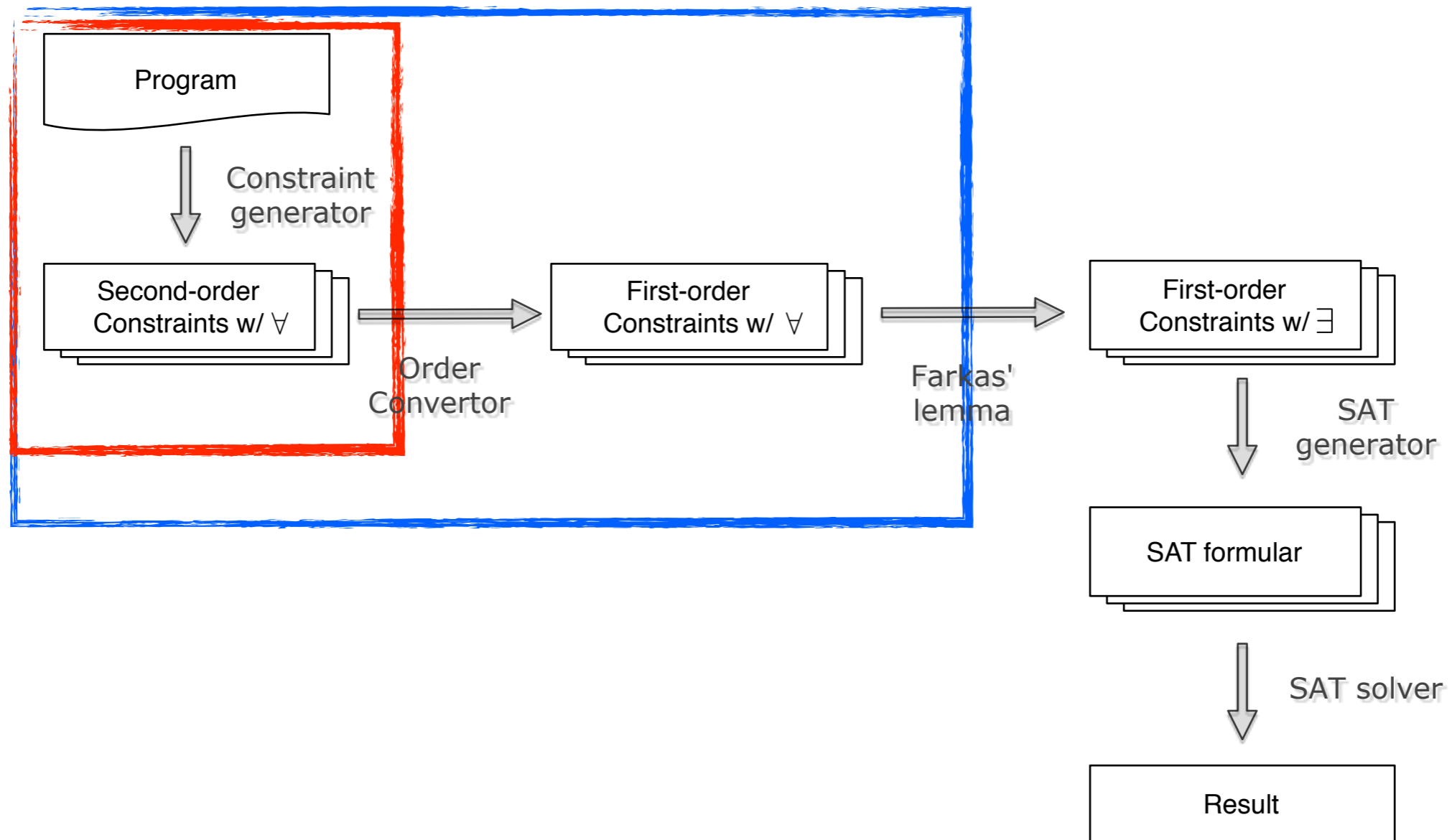
# Bird's-eye view



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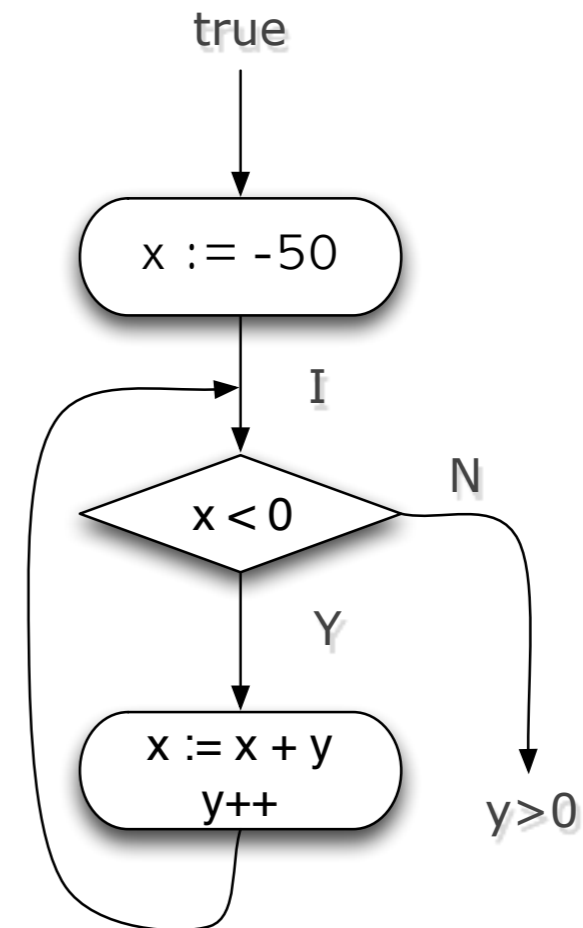


# Bird's-eye view



# Constraint generator

```
PV I (int y) {  
  x := -50;  
  while (x < 0) {  
    x := x + y;  
    y++;  
  }  
  assert(y > 0)  
}
```



$\forall x, y (I) :$

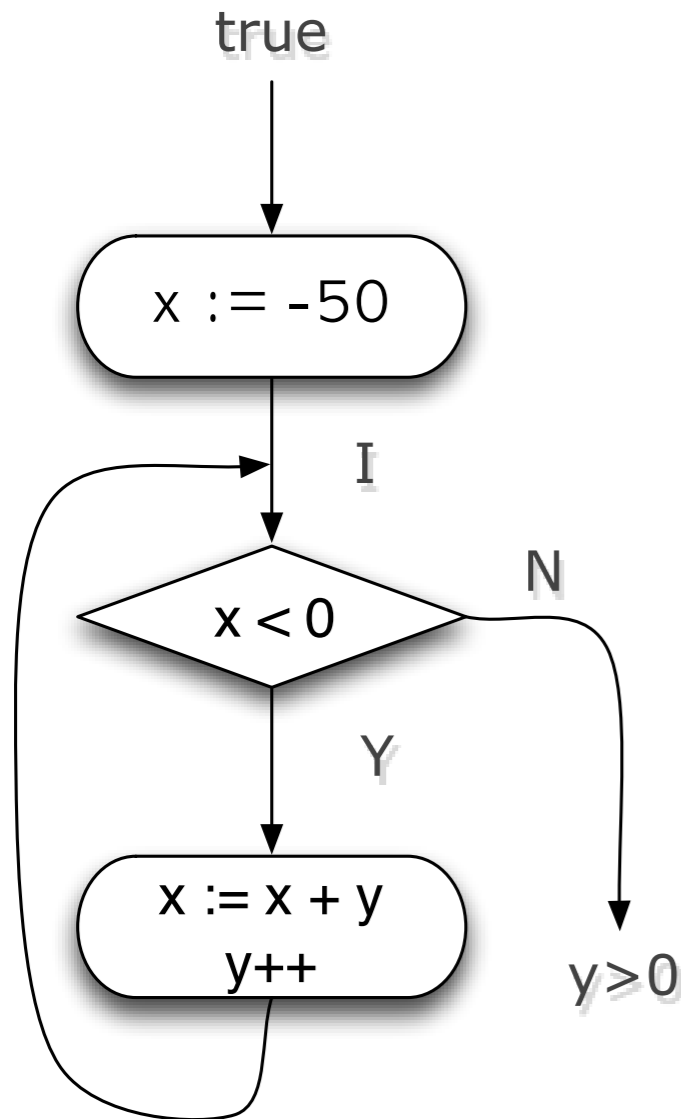
$\text{true} \Rightarrow I[-50/x]$

$I \wedge x < 0 \Rightarrow I[(y + 1)/y, (x + y)/x]$

$I \wedge x \geq 0 \Rightarrow y > 0$

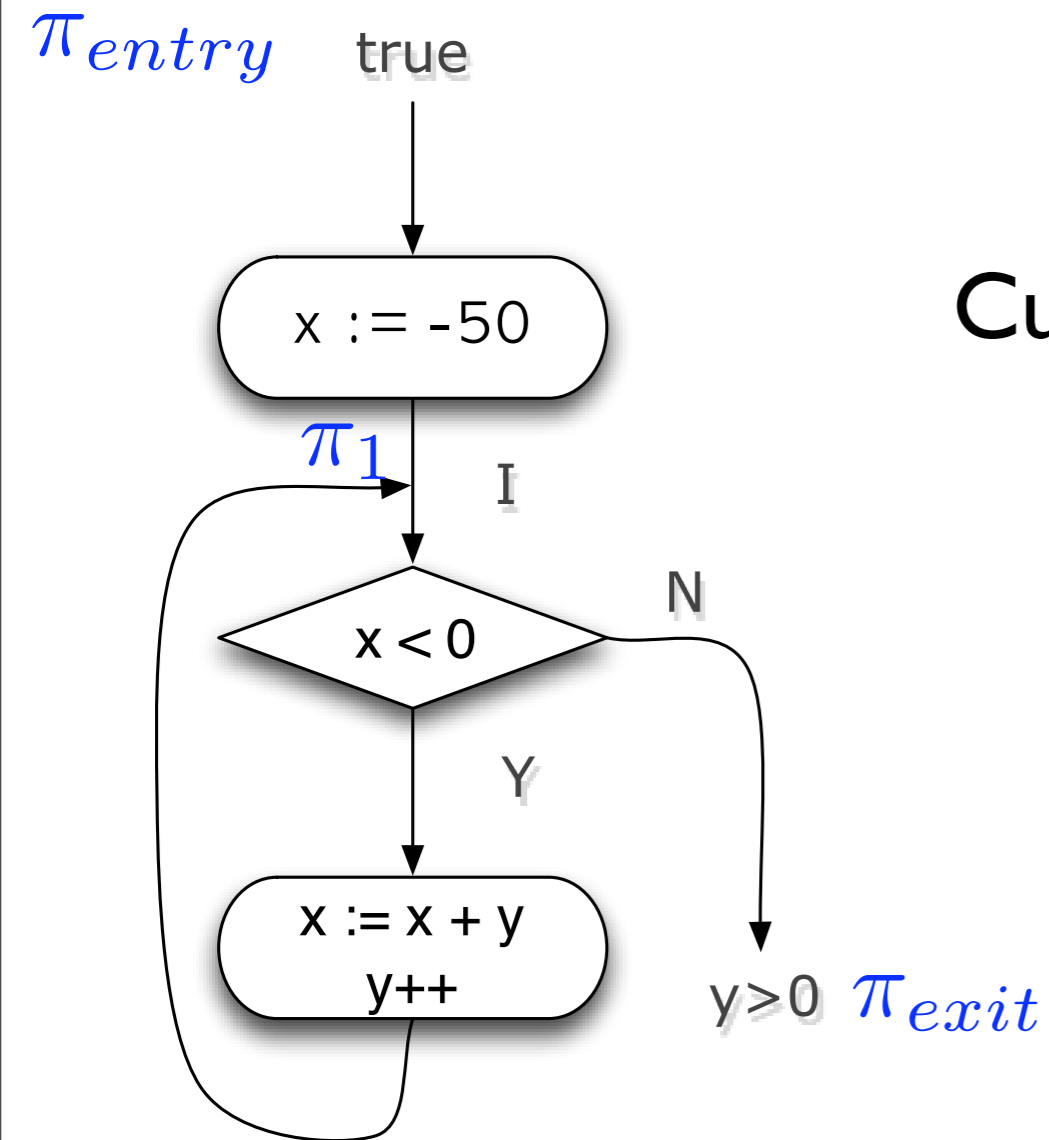
\*taken from the reference paper

# Constraint generator



Cut-set : {entry point, exit point, cut-points}

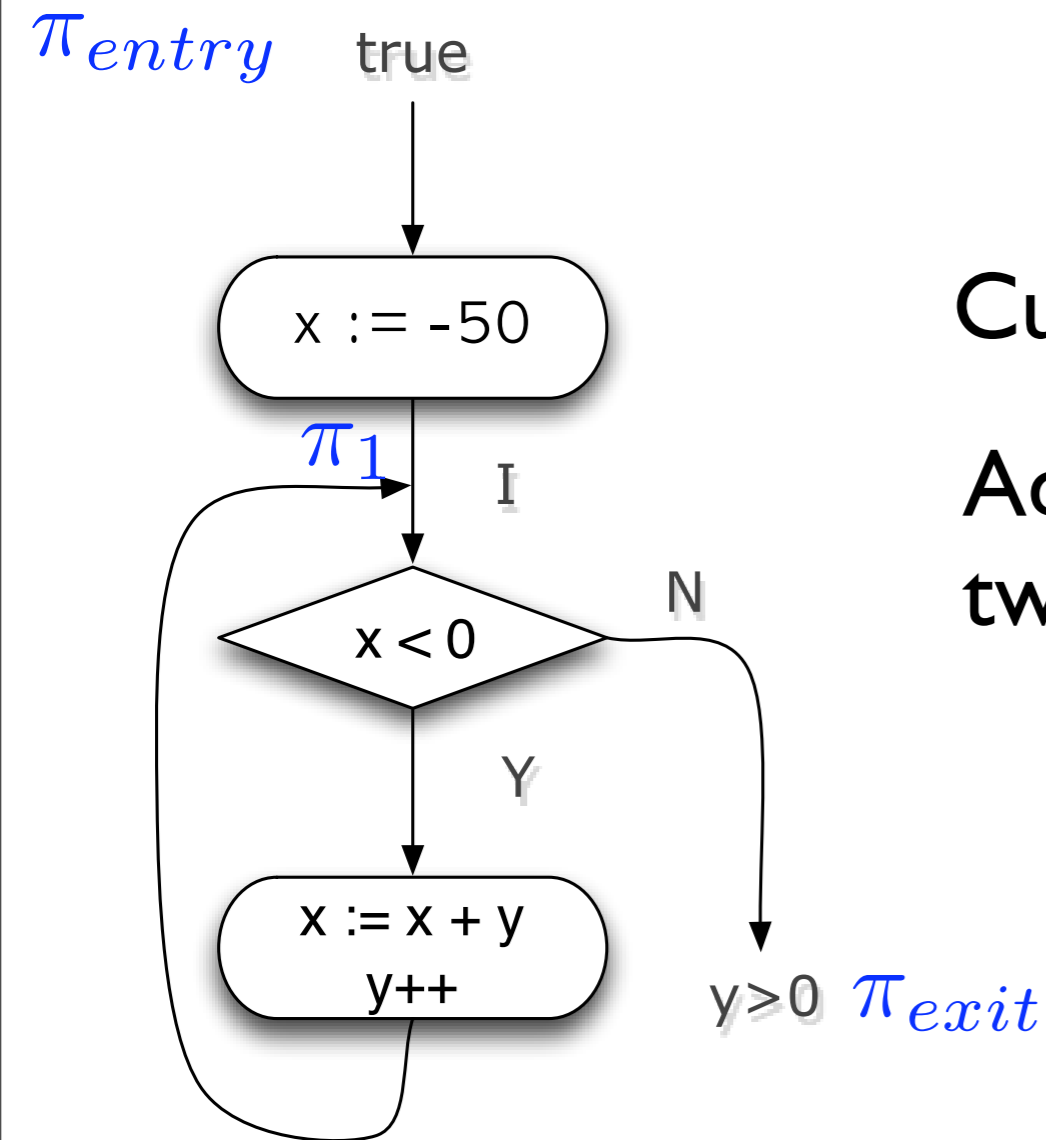
# Constraint generator



Cut-set :  $\{ \pi_{entry} , \pi_{exit} , \pi_1 \}$



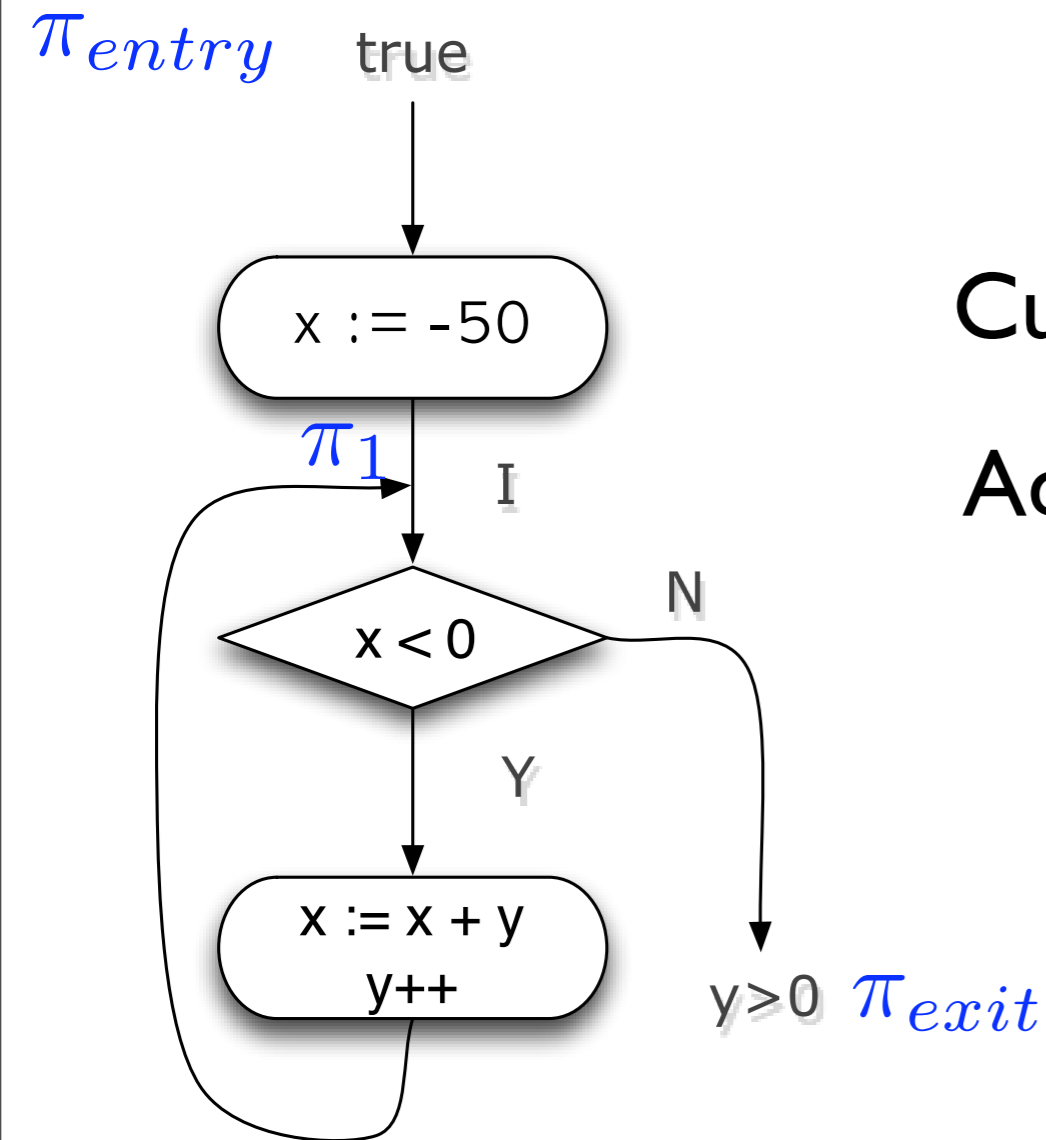
# Constraint generator



Cut-set :  $\{ \pi_{entry} , \pi_{exit} , \pi_1 \}$

Adjacent cut-points :  
two cut-points directly connected on a path

# Constraint generator

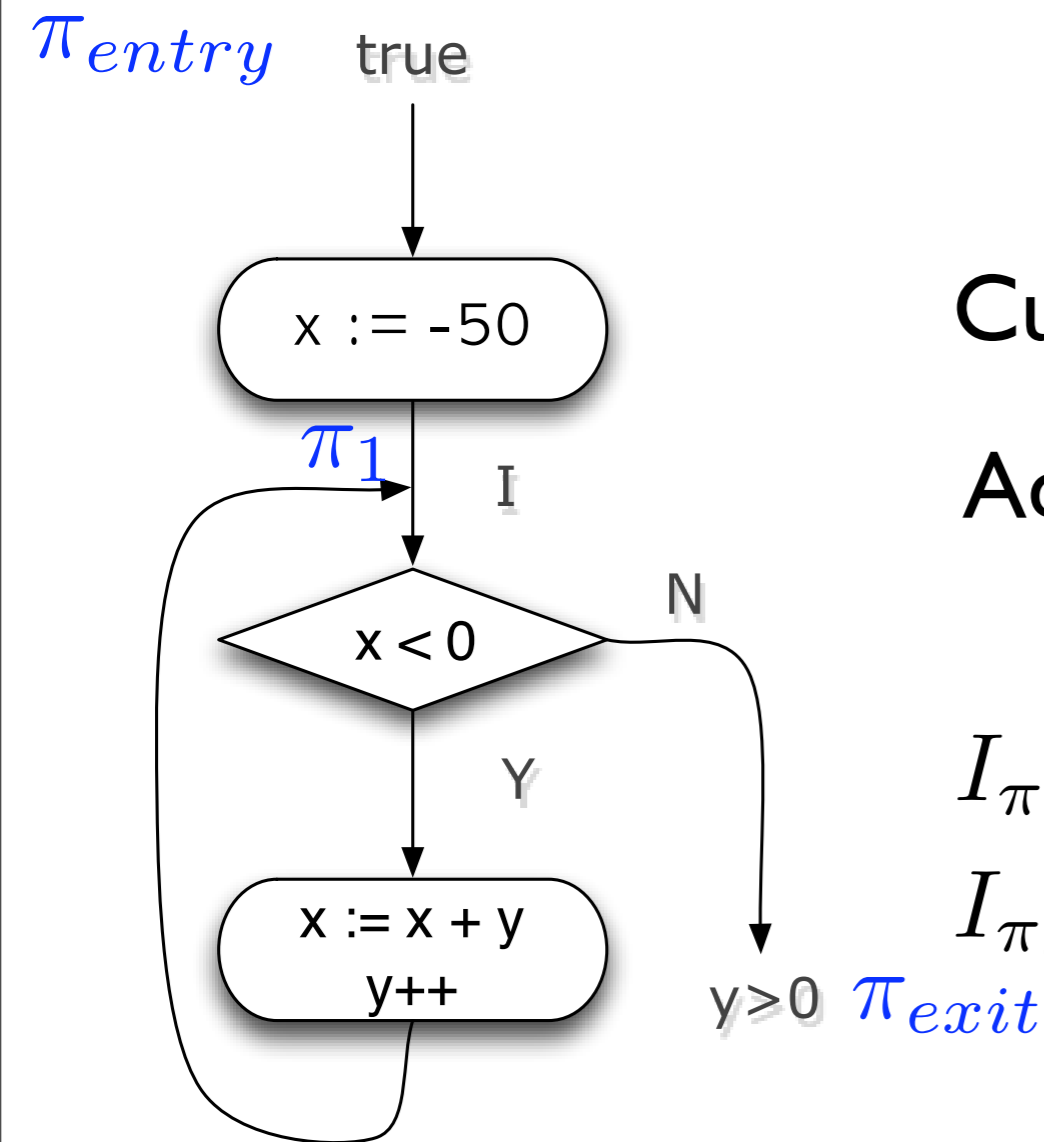


Cut-set :  $\{ \pi_{entry} , \pi_{exit} , \pi_1 \}$

Adjacent cut-points :

$(\pi_{entry}, \pi_1), (\pi_1, \pi_1), (\pi_1, \pi_{exit})$

# Constraint generator



Cut-set :  $\{ \pi_{entry} , \pi_{exit} , \pi_1 \}$

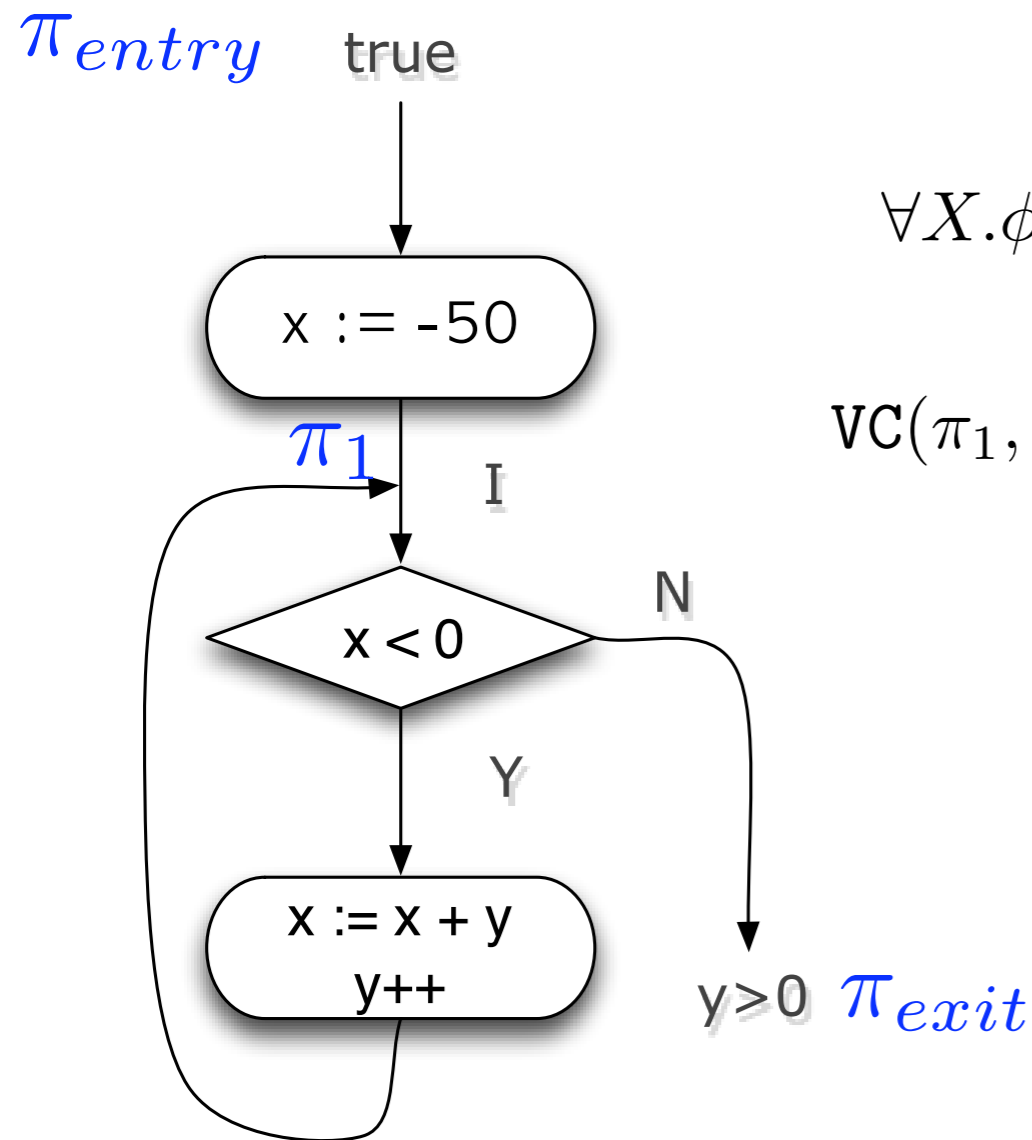
Adjacent cut-points :

$(\pi_{entry}, \pi_1), (\pi_1, \pi_1), (\pi_1, \pi_{exit})$

$I_\pi$ : a relation over variables that are live at  $\pi$

$I_{\pi_{entry}} = \text{true}$        $I_{\pi_{exit}} = \text{true}$

# Constraint generator



$$\forall X. \phi(I) = \bigwedge_{(\pi_1, \pi_2) \in \text{Adjacent}} VC(\pi_1, \pi_2)$$

$$VC(\pi_1, \pi_2) = \forall X \left( \bigwedge_{\tau \in \text{Paths}(\pi_1, \pi_2)} (I_{\pi_1} \Rightarrow \omega(\tau, I_{\pi_2})) \right)$$

$$\omega(\text{skip}, I) = I$$

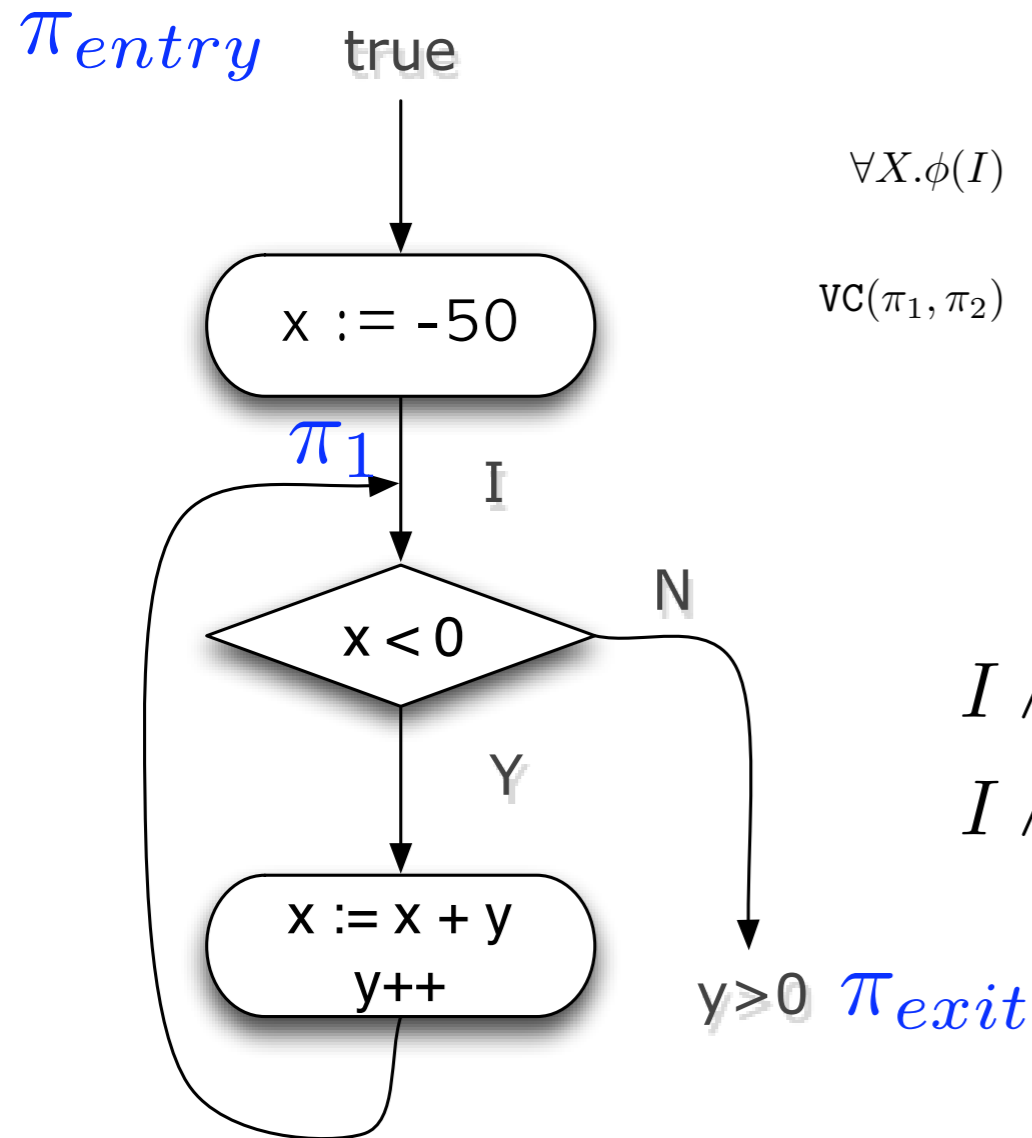
$$\omega(x := e, I) = I[e/x]$$

$$\omega(\text{assume } p, I) = p \Rightarrow I$$

$$\omega(\text{assert } p, I) = p \wedge I$$

$$\omega(S_1; S_2, I) = \omega(S_1, \omega(S_2, I))$$

# Constraint generator



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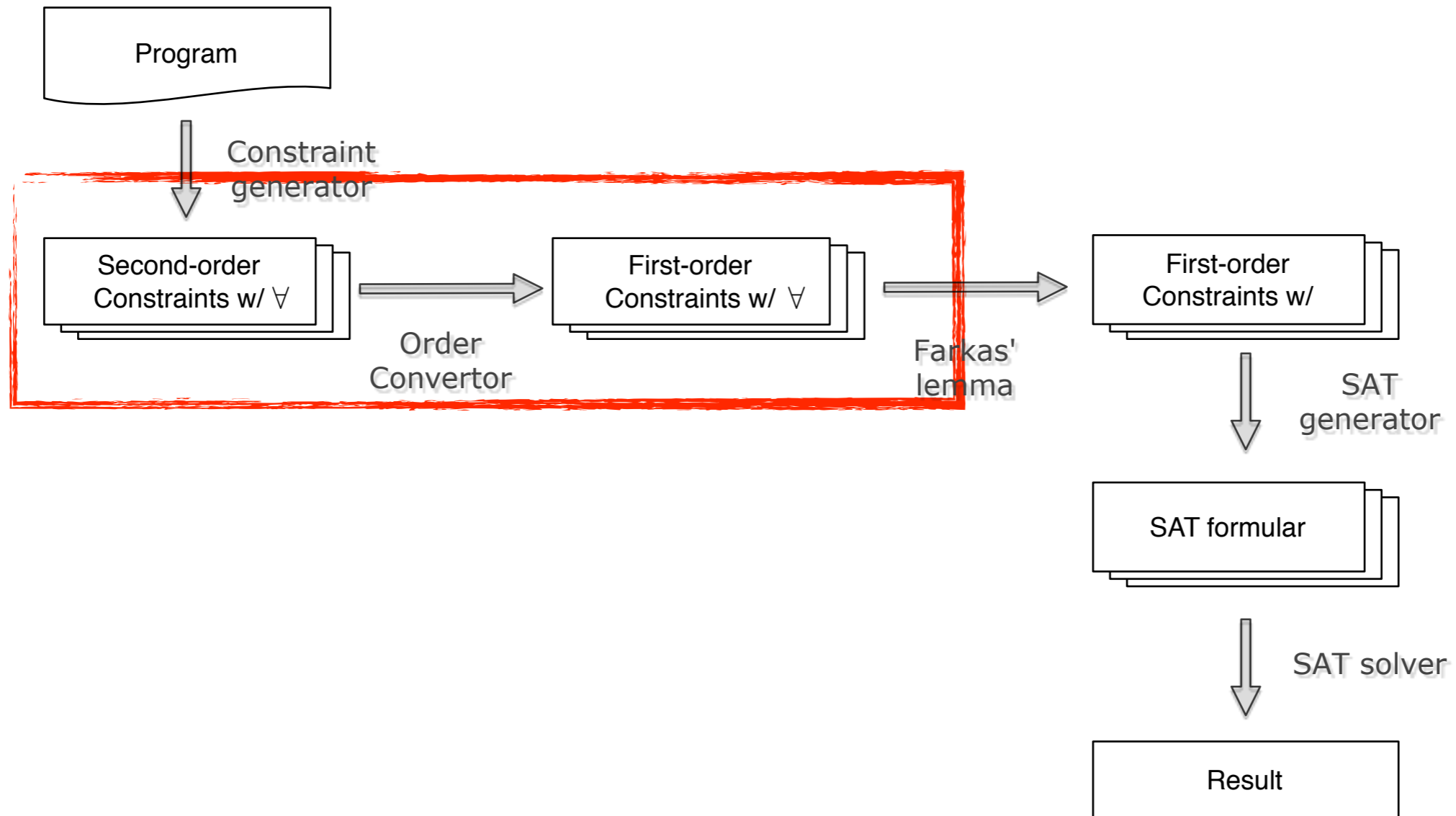
$$\omega(S_1; S_2, I) = \omega(S_1, \omega(S_2, I))$$

$$\text{true} \Rightarrow I[-50/x]$$

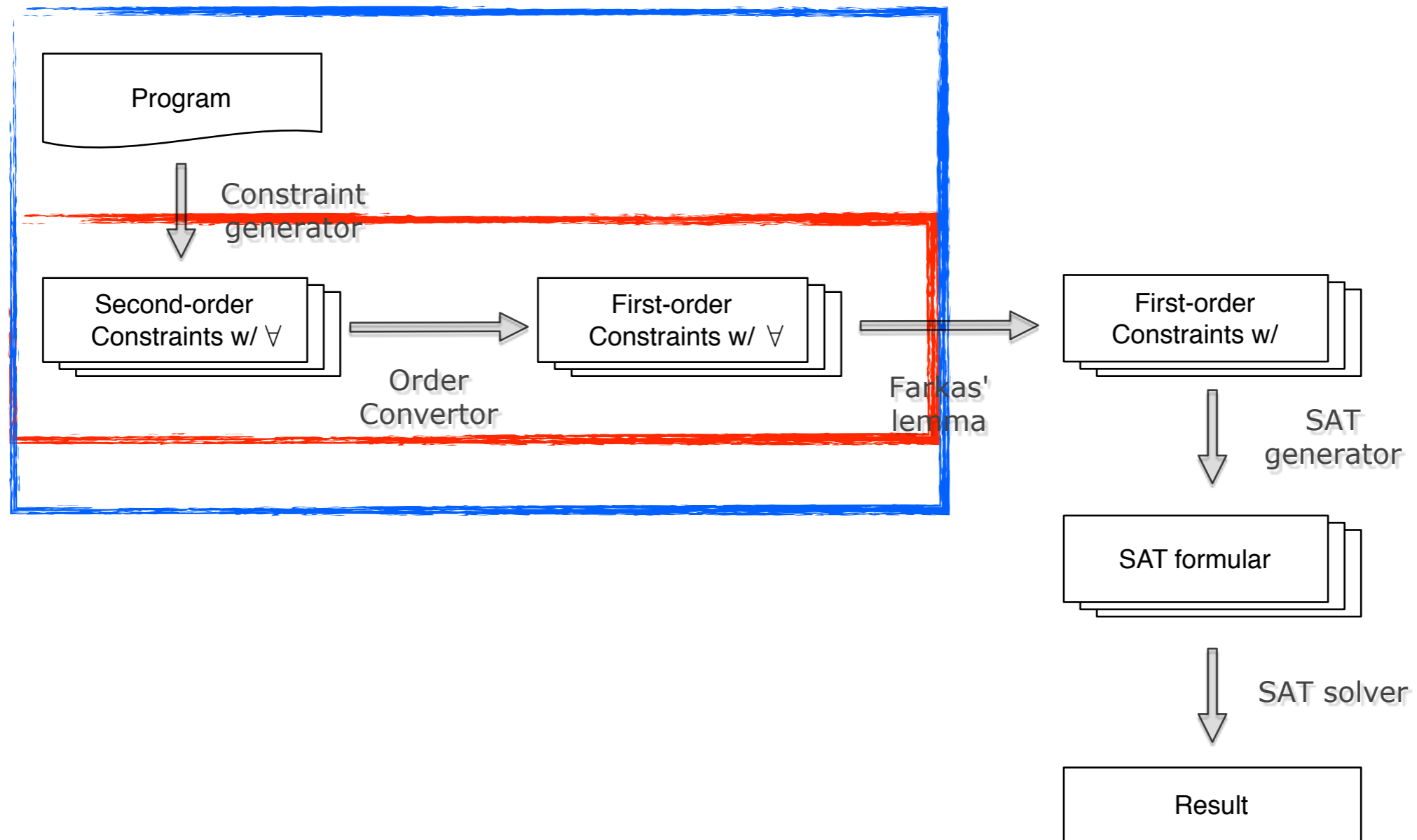
$$I \wedge x < 0 \Rightarrow I[(y+1)/y, (x+y)/x]$$

$$I \wedge x \geq 0 \Rightarrow y > 0$$

# Bird's-eye view



# Bird's-eye view



# Order convertor

$\forall_{x,y}(I) :$

$\text{true} \Rightarrow I[-50/x]$

$I \wedge x < 0 \Rightarrow I[(y + 1)/y, (x + y)/x]$

$I \wedge x \geq 0 \Rightarrow y > 0$

- Convert second-order constraint to first-order



# Order convertor

- Assume I has some form, e.g.,  $\sum_j a_j x_j \geq 0$
- Make some invariant templates
- Enumerate templates and find a right first-order constraint

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**Similar to our approach**

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For example, 3 kinds of templates

$$a_1x + a_2y + a_3 \geq 0$$

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$$\forall a_4x + a_5y + a_6 \geq 0$$

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$$\forall a_7x + a_8y + a_9 \geq 0$$

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Assume I :

$$\begin{aligned} a_1x + a_2y + a_3 &\geq 0 \\ \forall a_4x + a_5y + a_6 &\geq 0 \end{aligned}$$

$\forall_{x,y}\phi(a_j) :$

$$\text{true} \Rightarrow -50a_1 + a_2y + a_3$$

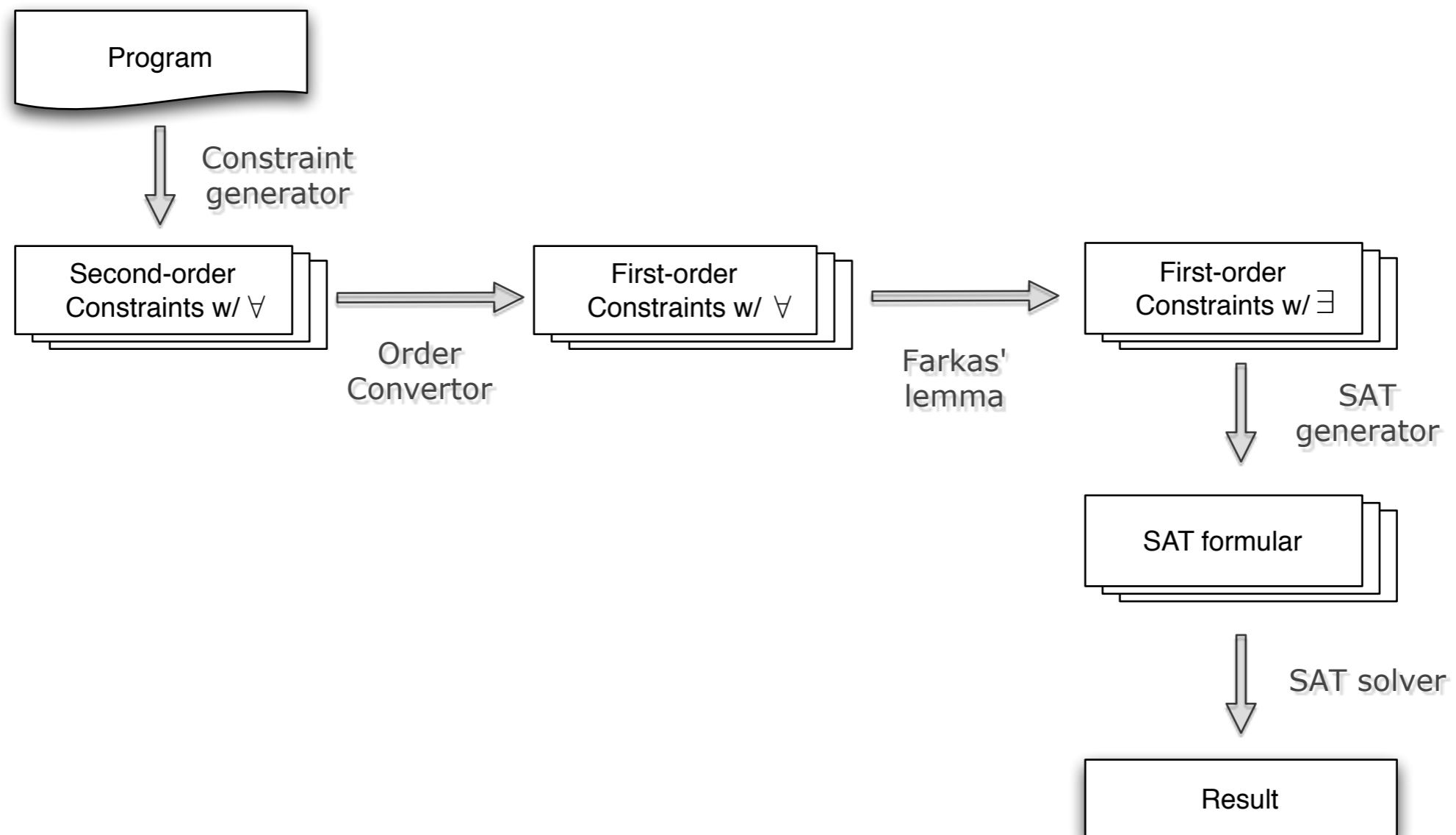
$$\vee -50a_4 + a_5y + a_6$$

$$I \wedge x < 0 \Rightarrow a_1(x+y+1) + a_2(y+1) + a_3$$

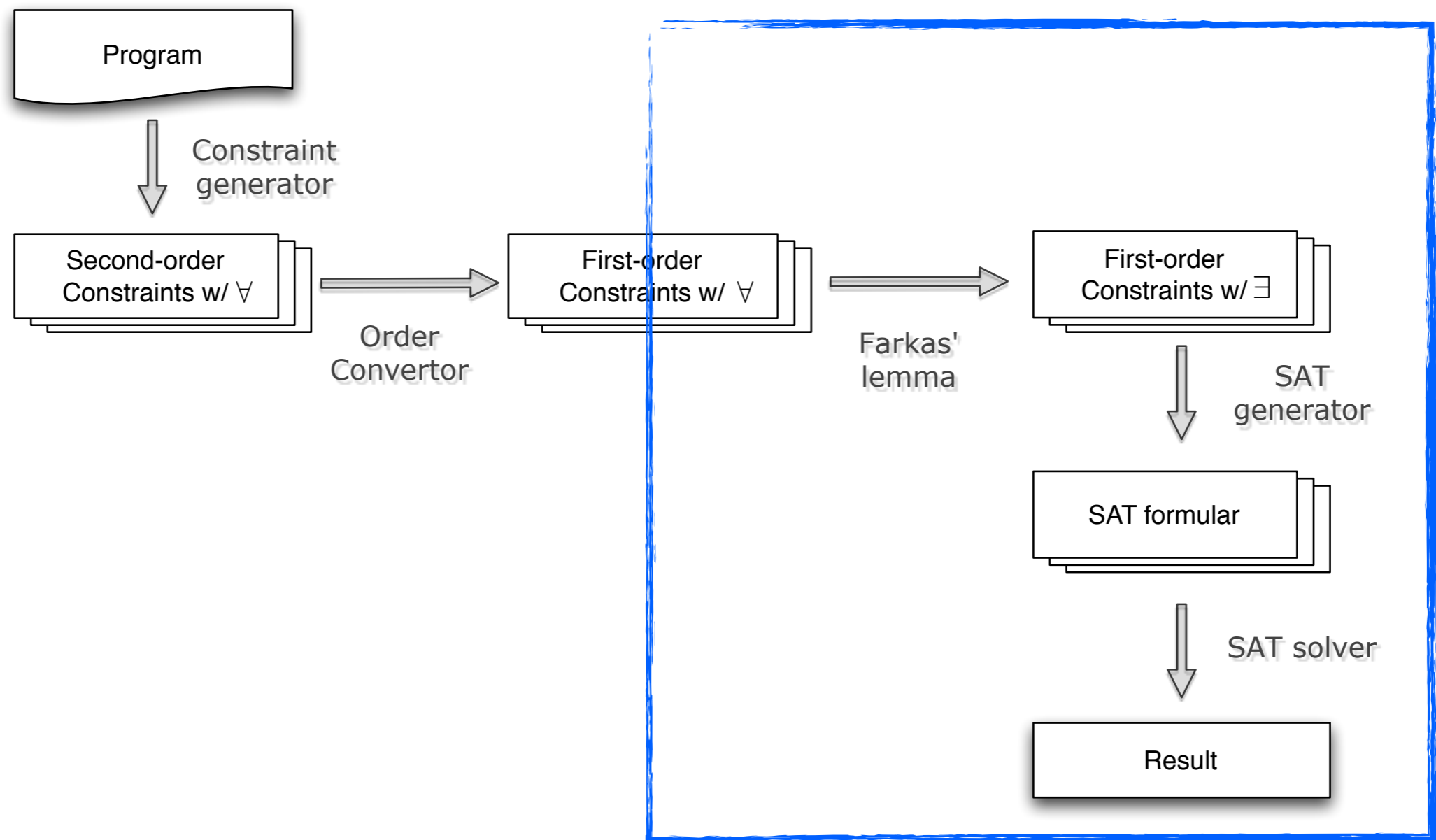
$$\vee a_4(x+y+1) + a_5(y+1) + a_6$$

$$I \wedge x \geq 0 \Rightarrow y > 0$$

# Bird's-eye view



# Bird's-eye view



# Experiment result

Name	This paper	BLAST	SLAM	GR06	ARMC
barbr	0.41	!	43.9	10.5	674
berkeley	3.00	!	2.90	0.10	4
bk-nat	5.30	!	!	0.43	3.25
seesaw	3.23	!	1.0	0.82	>2000

! : inability of tool to discover new predicates



# Conclusion

- **Very simple but powerful**
  - don't need fix-point calculation
- **Real-world invariants : Not that complicate**
  - Real-world data structure : Not that complicate (in our case)

Thank you