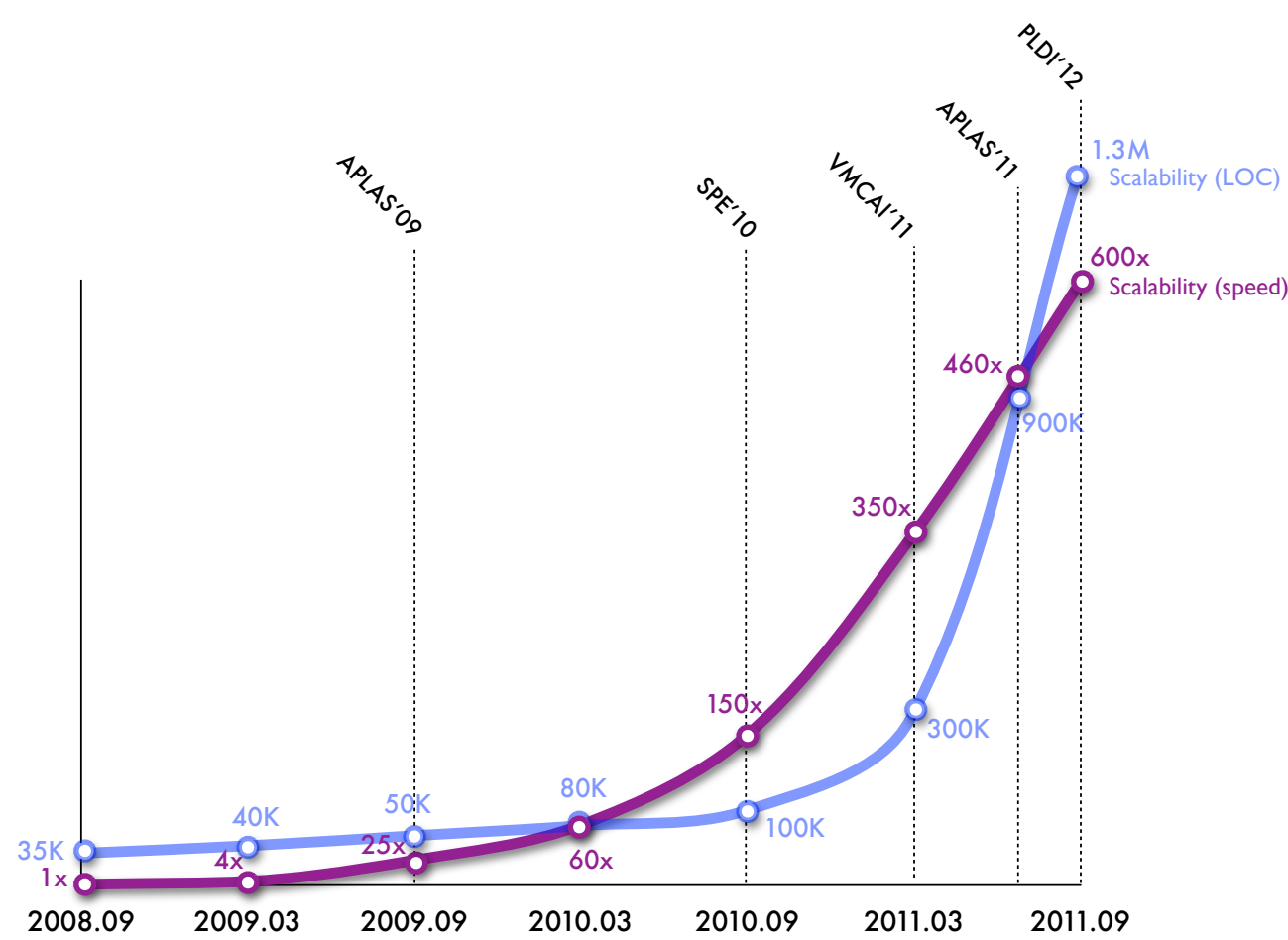


# ZooBerry: Automatic Generation of Sparse Global Static Analyzers & Their Validators

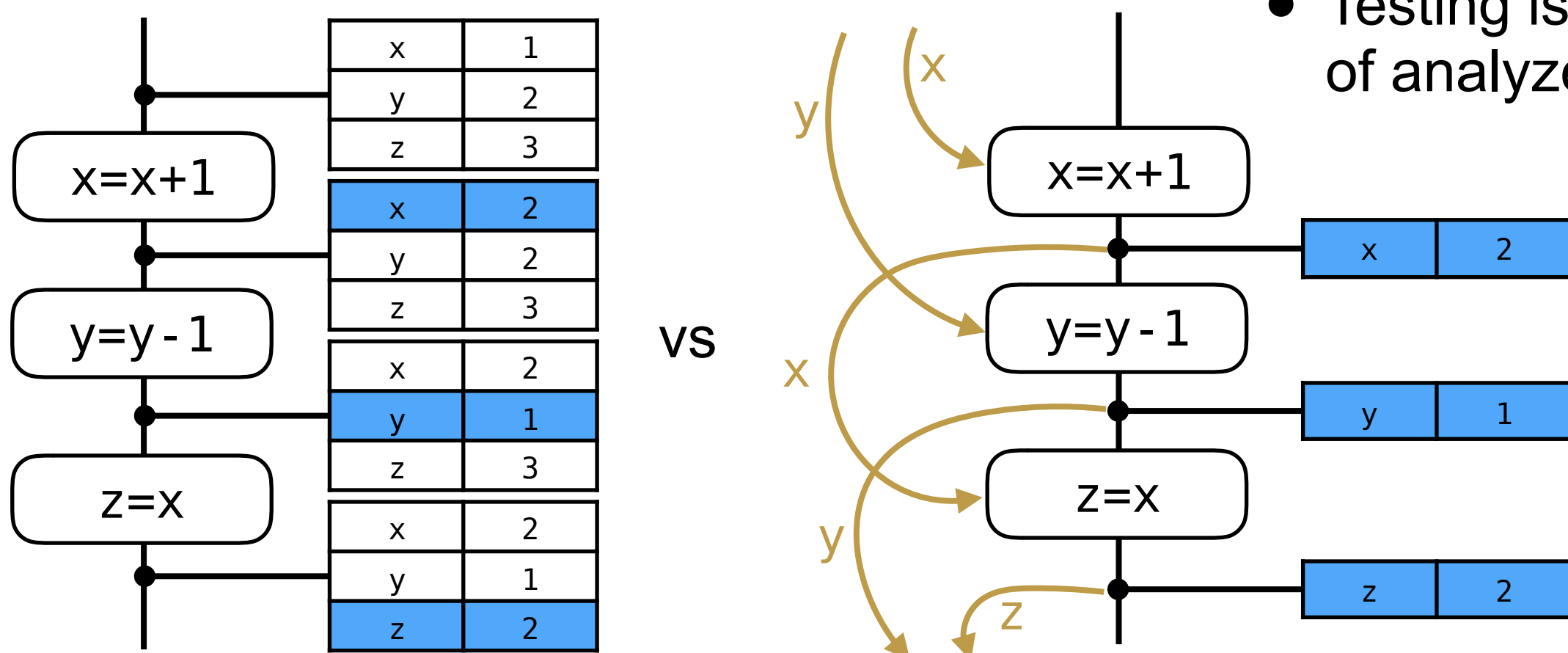
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## Background



### Key: sparse analysis<sup>1</sup>



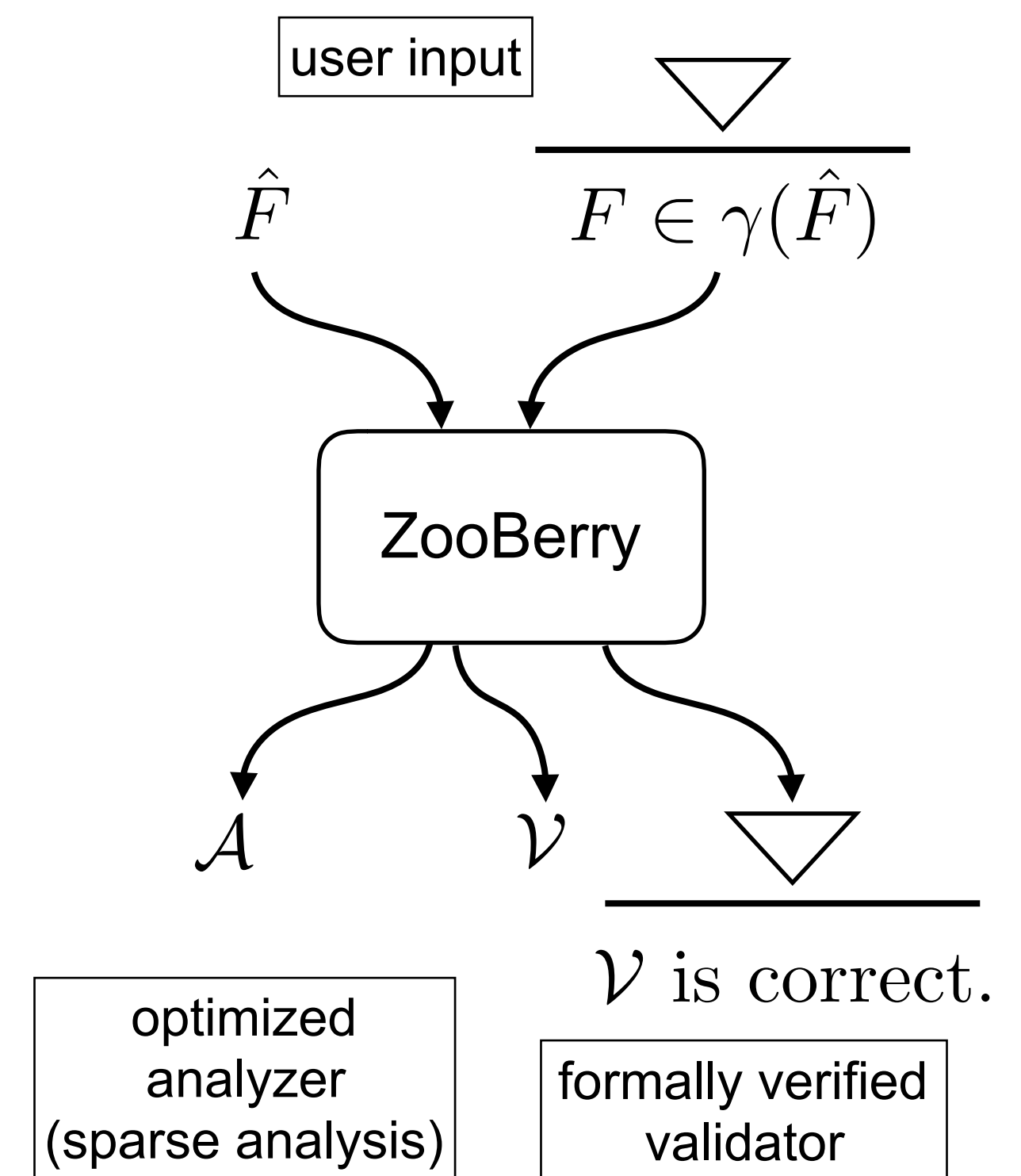
## Problem

Applying the sparse analysis is,

- burdensome
  - pre-analysis for drawing data dependency graph
  - additional 1500 LOC in OCaml (SparseSparrow<sup>1</sup>: our buffer overrun analyzer for C)
- (prone to be) buggy
  - 5 of 13 bugs we found<sup>2</sup> from SparseSparrow were about the sparse analysis.
  - Testing is not good to find bugs of analyzers.

## Solution

Automatic Generation!



## Automatically implementing the verified validation technique

### Correctness (naive)

A validation is simply a fixpoint check.

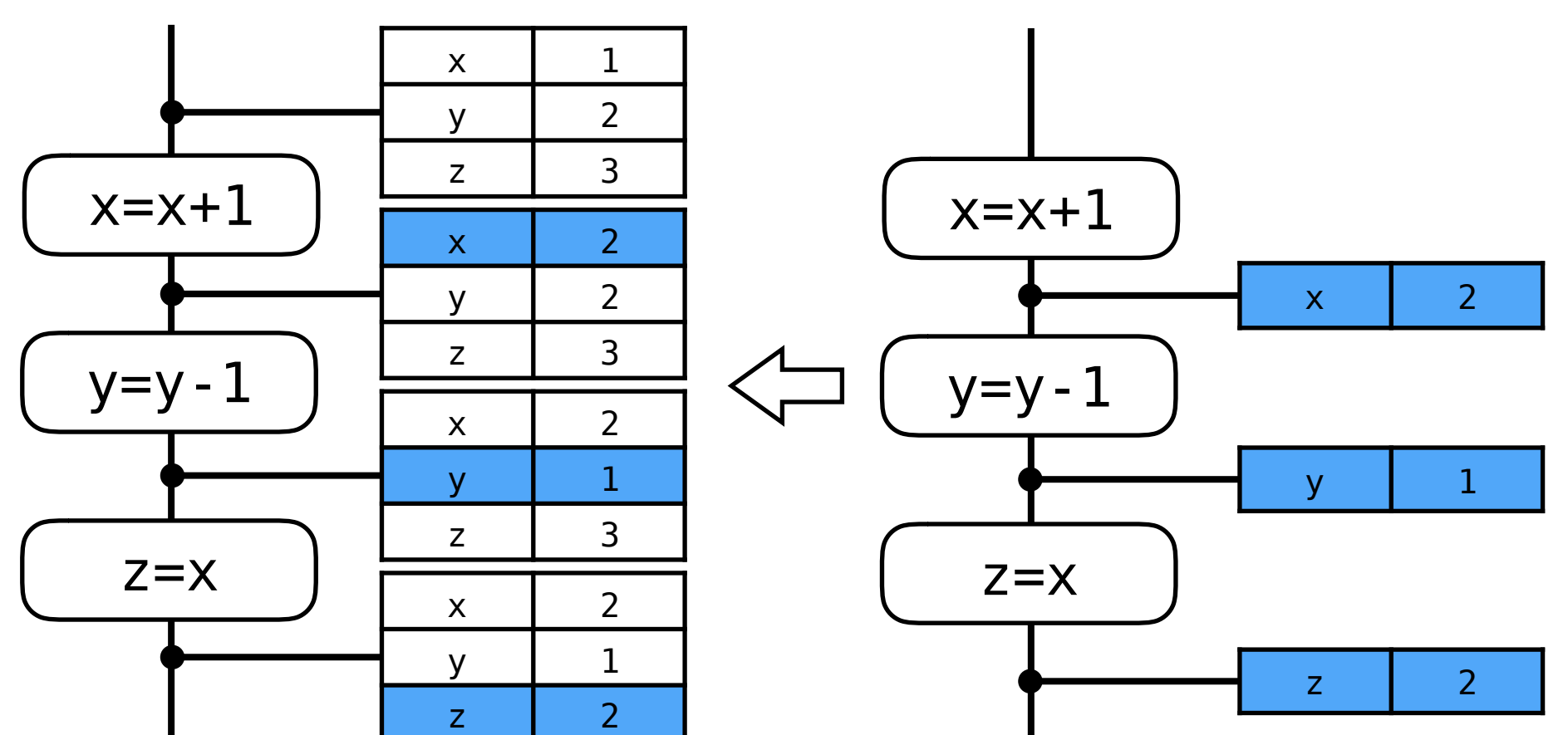
$$\forall p, \hat{s} : \mathcal{V}(p, \hat{s}) = \text{success} \Rightarrow \llbracket p \rrbracket \in \gamma(\hat{s})$$

$$\hat{F}(\hat{s}) \subseteq \hat{s}$$

However,  $\hat{f}(\text{y=y-1}, \llbracket x=2, y=2 \rrbracket) \not\subseteq \llbracket y=1 \rrbracket$

### Densification

"Recover the original analysis result and validate it."



- We designed an efficient densification algorithm.
- The validator checks also the densification is correct.

### Correctness (revised)

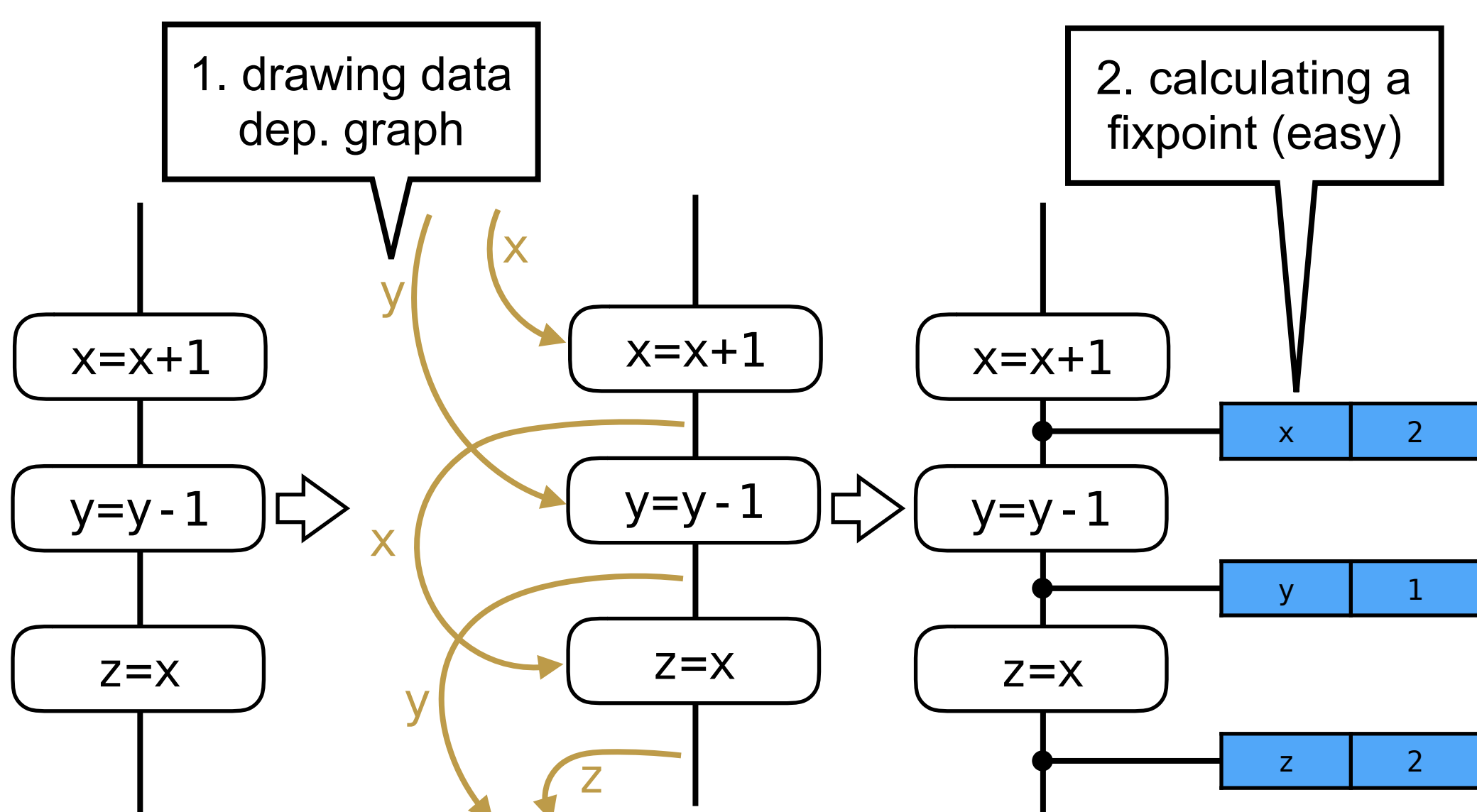
$$\forall p, \hat{s}, \mathcal{D}(\hat{s}) : \mathcal{V}(p, \hat{s}, \mathcal{D}(\hat{s})) = \text{success} \Rightarrow \llbracket p \rrbracket \in \gamma(\mathcal{D}(\hat{s})) \wedge \hat{s} \subseteq_{\text{use}}^p \mathcal{D}(\hat{s})$$

## References

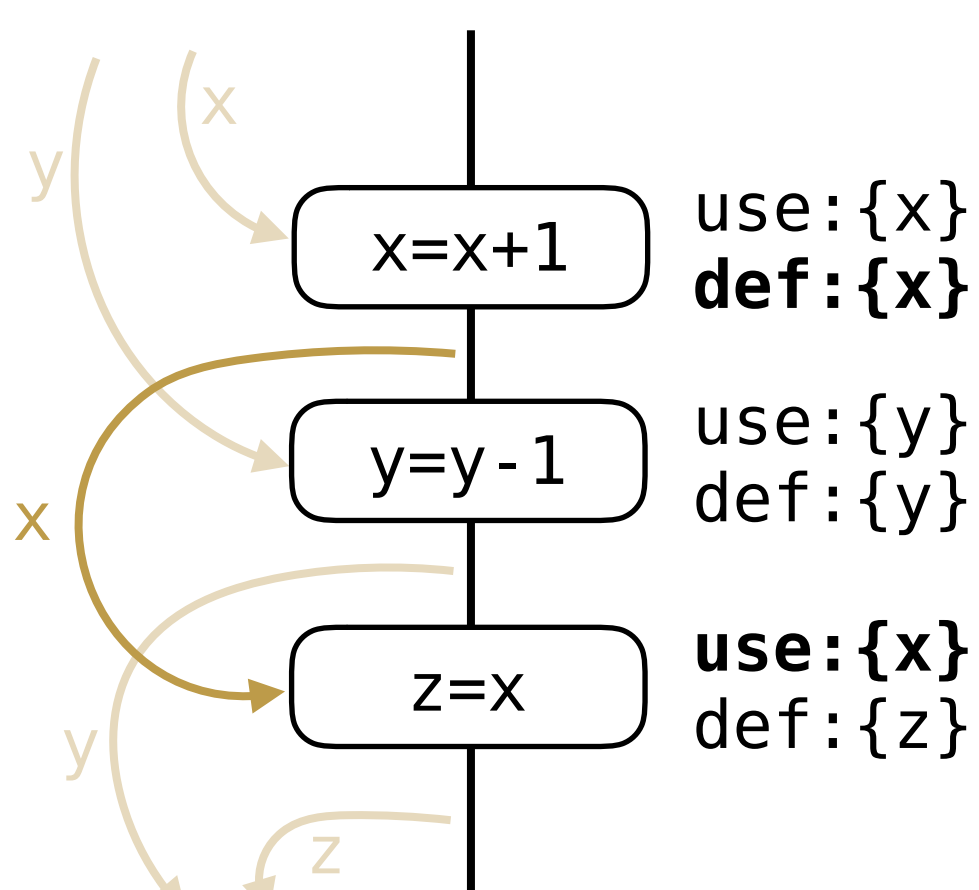
- Design and Implementation of Sparse Global Analyses for C-like Languages. Hakjoo Oh, Kihong Heo, Wonchan Lee, Woosuk Lee, and Kwangkeun Yi. *PLDI'12*.
- Towards Scalable Translation Validation of Static Analyzers. Jeehoon Kang, Sungkeun Cho, Joonwon Choi, Chung-Kil Hur, and Kwangkeun Yi. *ROSAEC techmemo 2014*. <http://rosaec.snu.ac.kr/publish/2014/techmemo/ROSAEC-2014-003.pdf>

## Automatically implementing the sparse analysis technique

### Applying the sparse analysis



access information  
set of locations that are defined/used during analyses



For sparsity, we designed a type-based transformation of the analysis spec.

$$\hat{f} : \text{Cmd} \rightarrow \hat{M} \rightarrow \hat{M}$$

$$\Downarrow$$

$$\hat{f}' : \text{Cmd} \rightarrow \hat{M} \rightarrow (\hat{M} \times \text{Acc})$$

$$\text{Acc} = 2^{\hat{L}} \times 2^{\hat{L}}$$