

4190.310 Programming Language Semeantics of K--

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$$\begin{array}{c}
 \text{TRUE} \frac{}{\sigma, M \vdash \mathbf{true} \Rightarrow \mathit{true}, M} \qquad \text{FALSE} \frac{}{\sigma, M \vdash \mathbf{false} \Rightarrow \mathit{false}, M} \\
 \\
 \text{NUM} \frac{}{\sigma, M \vdash \mathbf{n} \Rightarrow n, M} \qquad \text{UNIT} \frac{}{\sigma, M \vdash \mathbf{unit} \Rightarrow \cdot, M} \\
 \\
 \text{VAR} \frac{}{\sigma, M \vdash x \Rightarrow M(\sigma(x)), M} \\
 \\
 \text{ADD} \frac{\sigma, M \vdash E_1 \Rightarrow n_1, M' \quad \sigma, M' \vdash E_2 \Rightarrow n_2, M''}{\sigma, M \vdash E_1 + E_2 \Rightarrow n_1 + n_2, M''} \\
 \\
 \text{SUB} \frac{\sigma, M \vdash E_1 \Rightarrow n_1, M' \quad \sigma, M' \vdash E_2 \Rightarrow n_2, M''}{\sigma, M \vdash E_1 - E_2 \Rightarrow n_1 - n_2, M''} \\
 \\
 \text{MUL} \frac{\sigma, M \vdash E_1 \Rightarrow n_1, M' \quad \sigma, M' \vdash E_2 \Rightarrow n_2, M''}{\sigma, M \vdash E_1 \times E_2 \Rightarrow n_1 \times n_2, M''} \\
 \\
 \text{DIV} \frac{\sigma, M \vdash E_1 \Rightarrow n_1, M' \quad \sigma, M' \vdash E_2 \Rightarrow n_2, M''}{\sigma, M \vdash E_1 \div E_2 \Rightarrow n_1 \div n_2, M''} \\
 \\
 \text{EQUAL} \frac{\sigma, M \vdash E_1 \Rightarrow n_1, M' \quad \sigma, M' \vdash E_2 \Rightarrow n_2, M''}{\sigma, M \vdash E_1 = E_2 \Rightarrow n_1 = n_2, M''} \\
 \\
 \text{LESS} \frac{\sigma, M \vdash E_1 \Rightarrow n_1, M' \quad \sigma, M' \vdash E_2 \Rightarrow n_2, M''}{\sigma, M \vdash E_1 < E_2 \Rightarrow n_1 < n_2, M''} \\
 \\
 \text{NOT} \frac{\sigma, M \vdash E \Rightarrow b, M'}{\sigma, M \vdash \mathbf{not} E \Rightarrow \mathit{not} b, M'} \\
 \\
 \text{ASSIGN} \frac{\sigma, M \vdash E \Rightarrow v, M'}{\sigma, M \vdash x := E \Rightarrow v, M' \{ \sigma(x) \mapsto v \}}
 \end{array}$$

$$\begin{array}{c}
\text{SEQ} \frac{\sigma, M \vdash E_1 \Rightarrow v_1, M' \quad \sigma, M' \vdash E_2 \Rightarrow v_2, M''}{\sigma, M \vdash E_1; E_2 \Rightarrow v_2, M''} \\
\text{IFT} \frac{\sigma, M \vdash E \Rightarrow \text{true}, M' \quad \sigma, M' \vdash E_1 \Rightarrow v, M''}{\sigma, M \vdash \text{if } E \text{ then } E_1 \text{ else } E_2 \Rightarrow v, M''} \\
\text{IFF} \frac{\sigma, M \vdash E \Rightarrow \text{false}, M' \quad \sigma, M' \vdash E_2 \Rightarrow v, M''}{\sigma, M \vdash \text{if } E \text{ then } E_1 \text{ else } E_2 \Rightarrow v, M''} \\
\text{WHILEF} \frac{\sigma, M \vdash E_1 \Rightarrow \text{false}, M'}{\sigma, M \vdash \text{while } E_1 \text{ do } E_2 \Rightarrow \cdot, M'} \\
\text{WHILET} \frac{\sigma, M \vdash E_1 \Rightarrow \text{true}, M' \quad \sigma, M_1 \vdash \text{while } E_1 \text{ do } E_2 \Rightarrow v_2, M_2}{\sigma, M \vdash \text{while } E_1 \text{ do } E_2 \Rightarrow v_2, M'} \\
\text{FOR} \frac{\sigma, M \vdash E_1 \Rightarrow n_1, M' \quad \sigma, M' \vdash E_2 \Rightarrow n_2, M'' \quad \sigma, M'' \{\sigma(x) \mapsto n_1 + 0\} \vdash C \Rightarrow v_0, M_0 \quad \vdots \quad \sigma, M_{n_2-n_1-1} \{\sigma(x) \mapsto n_1 + (n_2 - n_1)\} \vdash C \Rightarrow v_{n_2-n_1}, M_{n_2-n_1}}{\sigma, M \vdash \text{for } x := E_1 \text{ to } E_2 \text{ do } C \Rightarrow \cdot, M_{n_2-n_1}} \quad n_2 \geq n_1 \\
\text{LETV} \frac{\sigma, M \vdash E_1 \Rightarrow v, M' \quad \sigma \{x \mapsto l\}, M' \{l \mapsto v\} \vdash E_2 \Rightarrow v', M''}{\sigma, M \vdash \text{let } x := E_1 \text{ in } E_2 \Rightarrow v', M''} \quad l \notin \text{Dom } M' \\
\text{LETF} \frac{\sigma \{f \mapsto \langle x, E_1, \sigma \rangle\}, M \vdash E_2 \Rightarrow v, M'}{\sigma, M \vdash \text{let proc } f(x) \text{ in } E_2 \Rightarrow v, M'} \\
\text{CALLV} \frac{\sigma, M \vdash E \Rightarrow v, M' \quad \sigma' \{x \mapsto l\} \{f \mapsto \langle x, E', \sigma' \rangle\}, M' \{l \mapsto v\} \vdash E' \Rightarrow v', M''}{\sigma, M \vdash f(E) \Rightarrow v', M''} \quad \sigma(f) = \langle x, E', \sigma' \rangle \quad l \notin \text{Dom } M' \\
\text{CALLR} \frac{\sigma' \{x \mapsto \sigma(y)\} \{f \mapsto \langle x, E, \sigma' \rangle\}, M \vdash E \Rightarrow M'}{\sigma, M \vdash f \langle y \rangle \Rightarrow M'} \quad \sigma(f) = \langle x, E, \sigma' \rangle \\
\text{READ} \frac{}{\sigma, M \vdash \text{read } x \Rightarrow n, M \{\sigma(x) \mapsto n\}} \\
\text{WRITE} \frac{\sigma, M \vdash E \Rightarrow v, M'}{\sigma, M \vdash \text{write } E \Rightarrow v, M'}
\end{array}$$